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Information Thermodynamics and Halting Problem

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Abstract

The formulations of the *undecidability* of the Halting Problem assume that the computing process, being observed, the description of which is given on the input of the 'observing' Turing Machine, is, at any given moment, the exact copy of the computing process running in the observing machine itself (the Cantor diagonal argument). In this way an *infinite cycle* is created shielding what is to be possibly discovered - the possible infinite cycle in the observed computing process. By this type of our consideration and in the thermodynamic sense the equilibrium status of a certain thermodynamic system is described or, even created. This is a thermodynamic image of the Cantor diagonal method used for seeking a possible infinite cycle and which, as such, has the property of the Perpetuum Mobile - the structure of which is recognizable and therefore we can avoid it. Thus we can show that it is possible to recognize the infinite cycle as a certain *original* equilibrium, but with a 'step-aside' or a time delay in evaluating the trace of the observed computing process.

The trace is a record of the *sequence of configurations* of the observed Turing machine. These configurations can be simplified to their *common configuration types*, creating now a word of a *regular language*. Furthermore, the control unit of any Turing Machine is a *finite automaton*. Both these facts enable the *Pumping Lemma* in the observing Turing Machine to be usable. In compliance with the Pumping Lemma, we know (the observing Turing Machine knows) that certain common configuration types must be *periodically repeated* in the case of the *infinite length* of their regular language. This fact enables (in a finite time) us (the observing Turing Machine) to decide that the observed computing process has entered into an infinite cycle.

Considerations of the real sense of the Gibbs Paradox are used to illustrate the idea of the term 'step-aside' which is our main methodological tool for looking for the infinite cycle in a Turing computing process and which enables us to avoid the commonly used attempts to solve the Halting Problem.

Keywords: Heat and Information Entropy, Observation, Carnot Cycle, Information Channel, Turing Machine, Infinite Cycle

1. Introduction

The formulations of the *undecidability* of the Halting Problem assume that the computing process, being observed, the description of which is given on the input of the 'observing' Turing Machine, is the exact copy of the computing process running in the observing machine itself (the Cantor diagonal argument in the Minski's proof [18]). By this way the *Auto-Reference* or an *infinite cycle* in computing sense or the *Self-Observation* in information sense or an analogue of the *stationary (equilibrium) state* in thermodynamic sense is created, *shielding* now what is to be possibly discovered - the infinite cycle in the observed computing process - also for its normal input. This shield is the real result of the Cantor diagonal argument which has been used for solving the Halting Problem, but on the contrary, creates it [18]. This shield is, also, a certain image of the sought possible infinite cycle. This shield could be, in the thermodynamic point of view, ceased or ended whether the performance of the Perpetuum Mobile functionality was possible, which is not (when, e.g., the equation $x = x + 1$ would be solvable) or by an external activity or approach.

This situation is recognizable and as such is *decidable* and solvable in all cases of its realizations by a certain '*step-aside*'. For this, we use the previously studied [5, 6, 9, 11] congruence between a cyclical thermodynamic process represented by the Carnot Cycle and a repeatable information transfer represented by the Shannon Transfer Chain but we enrich this effort now by their another congruence with the computing process running in the Turing Machine. Considerations of Gibbs Paradox [7, 8, 11] are used to illustrate the main idea of the term '*step-aside*' which is our main methodological tool for looking for an infinite cycle in a Turing computing process and which enables us to avoid the traditional attempts of solving the Halting Problem. The gap in their formulations is due to that fact that they assume that the computing process observes itself by following itself in the same '*time-click*' of its activity. For this case the claim of the unsolvability of such a situation is, of course, valid. But with a *time delay* or the '*step-aside*' [or a memory (a storage)] considered it works with '*another*' data and '*is able to see*' on its own previous activity as the '*normal*' input data. The idea of the '*step-aside*' is based just on the result giving the solution to Gibbs Paradox and its information meaning [7, 8, 11, 12, 14], and enables us to be in compliance with the *II. Principle of Thermodynamics* during such a process.

Thus, we show that it is possible to recognize the infinite cycle, but with the *time delay* or a *staging* (instead of the time delay the *staging* is usable, lasting a longer time interval in each of its repetition) in evaluating the *trace* of the observed computing process. The trace is a message or a record, both about the input data and about the structure of the computing process being observed (the *listing*, the *cross-references* and the *memory dump* in the language of programmers). In this phase the *observing* Turing Machine (we ourselves) is raising the question: "Is there an infinite cycle?" Following the trace the observing machine gains the answer. In our case, the trace is a *recorded sequence of configurations* of the *observed* Turing Machine. These configurations can be simplified to their *general configuration types* which create now a word of a *regular language* [12, 14]. Furthermore, the control unit of any Turing Machine is a *finite automaton*. Both these facts enable the *Pumping Lemma* in the observing Turing Machine to be used. In

compliance with the Pumping Lemma, we know (the observing Turing Machine knows) that certain general configuration types must be *periodically repeated* in the case of the *infinite length* of their regular language. This fact enables us (the observing Turing Machine) to decide that the observed computing process has entered into an infinite cycle. This event is performed in a finite time and is, by this way, recognizable in the finite time, too. When the described method is used it becomes an *instance of observation*. By application to 'itself' it becomes a *higher instance* of observation, now observing the trace of its previous instances. Thus the method of staging of the observed process will be used again. This method is applicable to all computing processes and as such it represents a contribution to the *dead lock* indication problem.

This sequence of ideas differs from the Cantor diagonal argument. Mainly it is achieved by the physical, especially by the thermodynamic and information (structure) type of our consideration respecting the *II. Principle of Thermodynamics* as the very principal root for methodological approaches of all types, both excluding and also enriching the mere empty logic.

2. Notion of Auto-Reference

Paradoxical claims (paradoxes, *noetical paradoxes*, *contradictions*, *antinomia*) have two parts - both parts are true, but the truth of one part denies the truth of the second part.

They can arise by not respecting the *metalanguage (semantic) level* - which is the higher level of our thinking about problems and the *language (syntactic) level* - which is the lower level of formulations of our 'higher' thoughts. Also they arise by not respecting a *double-level organization and description of measuring* - by not respecting the need of a '*step-aside*' of the observer from the observed. And also they arise by not respecting various *time clicks* in time sequences. As for the latter case they are in a contradiction with the *causality principle*. The common feature for all these cases is the *Auto-Reference* construction which itself, solved by itself, always states the requirement for ceasing the *II. Principle of Thermodynamics* and all its equivalents [8-14].

Let us us introduce the **Russel's criterion** for removing paradoxes¹: **Within the flow of our thinking and speech we need and must distinguish between two levels of our thinking and expressing in order not to fall in a paradoxical claim by mutual mixing and changing them.**

These levels are the higher one, the metalanguage (semantic) level and the lower one, the language (syntactic) level. Being aware of the existence of these two levels we prevent ourselves from their mutual mixing and changing, we prevent ourselves from *application* our *metalanguage claims on themselves* but now on the language level or vice versa.

We must be aware that our claims about properties of considered objects are *created* on the higher level, rather richer both semantically and syntactically than the lower one on which we really express ourselves about these objects. The words and meanings of this lower (and 'narrower') level are common to both of them. Our speech is *formulated and performed* on the lower level describing here our 'higher' thoughts and on which the objects themselves have been

1 B. Russel, L. Whitehead, *Principia Mathematica*, 1910, 1912, 1913 and 1927.

described, defined yet too, of course from the higher level, but with the necessary (lower) limitations. (As such they are thought over on the higher level.) From this point of view we understand the various meanings (levels) of the same words. Then any mutual mixing and changing the metalanguage and language level or the autotereference (paradox, noetical paradox, contradiction, antinomium) is excluded.

2.1. Auto-Reference in Information Transfer, Self-Observation

In any *information transfer channel* \mathcal{K} the *channel equation*

$$H(X) - H(X|Y) = H(Y) - H(Y|X) \quad (1)$$

it is valid [23]. This equation describes the mutual relations among *information entropies* [(average) *information amounts*] in the channel \mathcal{K} .

The quantities $H(X)$, $H(Y)$, $H(X|Y)$ and $H(Y|X)$ are the *input*, the *output*, the *loss* and the *noise entropy*.

The difference $H(X) - H(X|Y)$ or the difference $H(Y) - H(Y|X)$ defines the *transinformation* $T(X;Y)$ or the transinformation $T(Y;X)$ respectively,

$$H(X) - H(X|Y) \triangleq T(X;Y) = T(Y;X) \triangleq H(Y) - H(Y|X) \quad (2)$$

When the channel \mathcal{K} transfers the information (entropy) $H(X)$, but now just at the value of the entropy $H(X|Y)$, $H(X) = H(X|Y)$, then, necessarily, must be valid

$$T(X;Y) = 0 \quad [= H(Y) - H(Y|X)] \quad (3)$$

- For $H(Y|X) = 0$, we have $T(X;Y) = H(Y) = 0$.
- For $H(Y|X) \neq 0$ we have $H(Y) = H(Y|X) \neq 0$

In both these two cases the channel \mathcal{K} operates as the *interrupted* (with the absolute noise) and the output $H(Y)$ is without any relation to the input $H(X)$ and, also, it doesn't relate to the structure of \mathcal{K} . This structure is expressed by the value of the quantity $H(X|Y)$. We assume, for simplicity, that $H(Y|X) = 0$.

From the (1)-(3) follows that the **channel** \mathcal{K} **can't** transfer (within the same step p of its *transfer process*) such an information which describes its inner structure and, thus, it can't **transfer - observe** (copy, measure) **itself**. It is valid both for the concrete information value and for the average information value, as well.

Any channel \mathcal{K} can't transfer its own states considered as the input messages (within the same steps p). Such an attempt is the information analogy for the **Auto-Reference** known from Logics and Computing Theory. Thus a certain 'step-aside' leading to a nonzero transfer output, $H(Y) = H(X) - H(X | Y) > 0$, is needed.

2.2. Auto-Reference and Thermodynamic Stationarity

The transfer process running in an information transfer channel \mathcal{K} is possible to be comprehended (modeled or, even, constructed) as the *direct* Carnot Cycle \mathcal{O} [6, 8]. The relation $\mathcal{O} \cong \mathcal{K}$ is postulated. Further, we can imagine its observing method, equivalent to its 'mirror' $\mathcal{O}'' \cong \mathcal{K}''$. This *mirror* \mathcal{O}'' is, at this case, the *direct* Carnot Cycle \mathcal{O} as for its structure, but functioning in the *indirect, reverse* mode [6, 8].

Let us connect them together to a *combined heat cycle* $\mathcal{O}\mathcal{O}''$ in such a way that the *mirror* (the *reverse cycle* \mathcal{O}'') is gaining the message about the structure of the *direct cycle* \mathcal{O} . This message is (carrying) the information $H(X | Y)$ about the structure of the transformation (transfer) process ($\mathcal{O} \cong \mathcal{K}$) being 'observed'. The *mirror* $\mathcal{O}'' \cong \mathcal{K}''$ is gaining this information $H(X | Y)$ on its noise 'input' $H(Y'' | X'')$ [while $H(X'') = H(Y)$ is its input entropy].

The quantities ΔQ_W , ΔA and ΔQ_0 or the quantities $\Delta Q''_W$, $\Delta A''$ and $\Delta Q''_0$ respectively, define the information entropies of the information transfer realized (thermodynamically) by the *direct* Carnot Cycle \mathcal{O} or by the *reverse* Carnot Cycle \mathcal{O}'' (the *mirror*) respectively (the *combined cycle* $\mathcal{O}\mathcal{O}''$ is created),

$$\begin{aligned} H(X) &= \frac{\Delta Q_W}{kT_W}, \text{ resp. } H(Y'') = \frac{\Delta Q''_W}{kT''_W} \\ H(Y) &= \frac{\Delta A}{kT_W}, \text{ resp. } H(X'') = \frac{\Delta A''}{kT''_W} \\ H(X|Y) &= \frac{\Delta Q_0}{kT_W}, \text{ resp. } H(Y''|X'') = \frac{\Delta Q''_0}{kT''_W} \end{aligned} \quad (4)$$

Our aim is to gain the *nonzero* output mechanical work ΔA^* of the combined heat cycle $\mathcal{O}\mathcal{O}''$, $\Delta A^* > 0$. We want to gain nonzero information $H^*(Y^*) = \Delta A^* / kT_W > 0$.

To achieve this aim, for the efficiencies η_{max} and η''_{max} of the both connected cycles \mathcal{O} and \mathcal{O}'' (with the working temperatures $T_W = T''_W$ and $T_0 = T''_0$, $T_W \geq T_0 > 0$), it must be valid that $\eta_{max} > \eta''_{max}$; we want the validity of the relation²

$$\Delta^* A = \Delta A - \Delta A'' > 0 \quad [\Delta A'' = \Delta Q''_W - \Delta Q''_0] \quad (5)$$

² We follow the proof of *physical and thus logical impossibility* of the construction and functionality of the Perpetuum Mobile of the *II.* and, equivalently [8], of the *I.* type.

When $\Delta Q_0 = \Delta Q''_0$ should be valid, then must be that $\Delta Q''_w < \Delta Q_W \Leftrightarrow (\eta_{max} > \eta''_{max})$ and thus it should be valid that

$$\begin{aligned} \Delta A^* &= \Delta Q_W \cdot \eta_{max} - \Delta Q''_W \cdot \eta''_{max} > 0 \text{ but} \\ \Delta Q_W \cdot \eta_{max} - \Delta Q''_W \cdot \eta''_{max} &= \Delta Q_0 - \Delta Q''_0 = 0 \end{aligned} \quad (6)$$

Thus the output work $\Delta A^* > 0$ should be generated without any lost heat and by the direct change of the whole heat $\Delta Q_W - \Delta Q''_w$ but within the cycle $\mathcal{O}\mathcal{O}''$. For $\eta_{max} < \eta''_{max}$ the same heat $\Delta Q_W - \Delta Q''_w$ should be pumped from the cooler with the temperature T_0 to the heater with the temperature T_W directly, without any compensation by a mechanical work. We see that $\Delta A^* = 0$ is the reality.

Our combined machine $\mathcal{O}\mathcal{O}''$ should be the *II. Perpetuum Mobile* in both two cases. Thus $\eta_{max} = \eta''_{max}$ must be valid (the heater with the temperature T_W and the cooler with the temperature T_0 are common) that

$$\eta_{max} = \eta''_{max} < 1 \text{ and then } \Delta Q_W = \Delta Q''_W \quad (7)$$

We must be aware that for $\eta_{max} = \eta''_{max} < 1$ the whole information entropy of the environment in which our (reversible) combined cycle $\mathcal{O}\mathcal{O}''$ is running changes on one hand by the value

$$H(X) \cdot \eta_{max} = \frac{\Delta Q_W}{kT_W} \cdot (1 - \beta) > 0, \quad \beta = 1 - \eta_{max} = \frac{T_0}{T_W} \quad (8)$$

and on the other hand it is also changed by the value $-H(X) \cdot \eta_{max} = -\frac{\Delta Q_W}{kT_W} \cdot (1 - \beta)$

Thus it must be changed by the zero value

$$H^*(Y^*) = \frac{\Delta A^*}{kT_W} = H(X) \cdot \eta_{max} - H(Y'') \cdot \eta''_{max} = H(X) \cdot (\eta_{max} - \eta''_{max}) = 0 \quad (9)$$

The whole combined machine, or the thermodynamic system with the cycle $\mathcal{O}\mathcal{O}''$ is, when the cycle $\mathcal{O}\mathcal{O}''$ is seen, as a whole, in the *thermodynamic equilibrium*. (It can be seen as an unit, analogous to an interruptable operation in computing.)

Thus, the observation of the observed process \mathcal{O} by the observing reverse process \mathcal{O}'' with the same structure (by itself), or the Self-Observation, is impossible in a physical sense, and, consequently, in a logical sense, too (see the Auto-Reference in computing).

Nevertheless, the construction of the Auto-Reference is describable and, as such, is recognizable, decidable just as a construction *sui generis*. It leads, necessarily, to the requirement of the *II. Perpetuum Mobile* functionality when the requirements (5) and (6) are sustained.

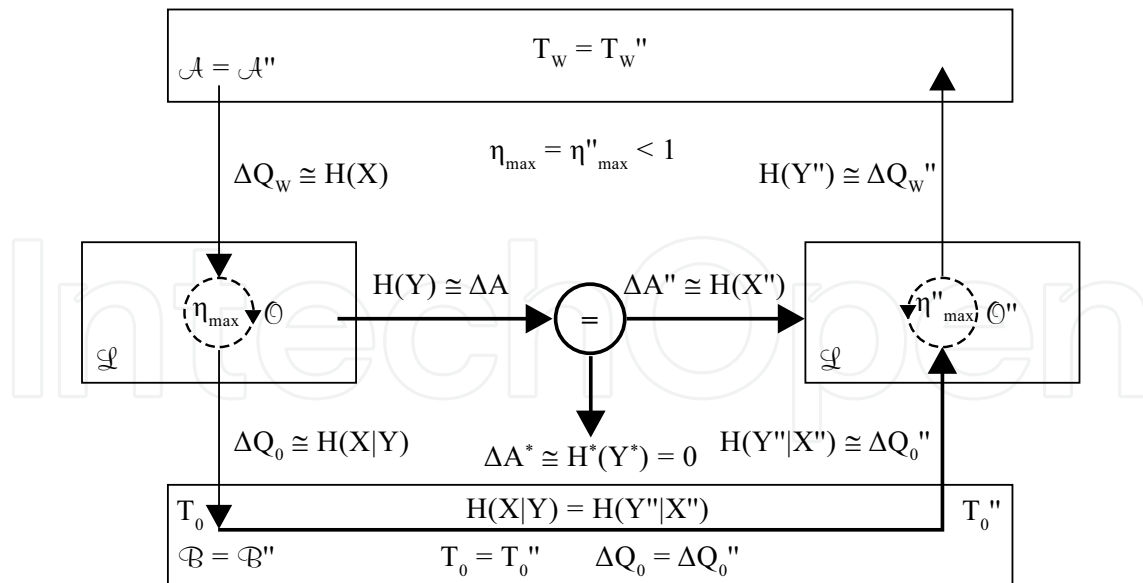


Figure 1. Stationarity of the double cycle $\mathcal{O}\mathcal{O}''$

(Note, that the Carnot Machine itself is, by its definition, a construction of the infinite cycle of the states of its working medium and as such is identifiable and recognizable.)

2.3. Gibbs Paradox and Auto-Reference in Observation and Information Transfer

Only just by a (thought) 'dividing' of an equilibrium system \mathcal{A} by *diaphragms* [20], without any influence on its thermodynamic (macroscopic) properties, a non-zero difference of its entropy, before and after its 'dividing', is evidenced.

Let us consider a thermodynamic system \mathcal{A} in volume V and with n matter units of ideal gas in the thermodynamic equilibrium. The *state equation* of \mathcal{A} is $pV = nR\Theta$. For an elementary change of the *internal energy* U of \mathcal{A} we have $dU = nc_v d\Theta$.

From the state equation of \mathcal{A} , and from the general *law of energy conservation* [for a (substitute) reversible exchange of heat δq between the system and its environment] we formulate the *I. Principle of Thermodynamics*, $\delta q = dU + pdV$

From this principle, and from *Clausius equation* $\Delta S \stackrel{\text{Def}}{=} \frac{\Delta q}{\Theta}, \Delta q = c_v \Delta\Theta + \frac{R\Theta \Delta V}{V}, \Theta > 0$, follows that

$$S = n \int \left(c_v \frac{d\Theta}{\Theta} + R \frac{dV}{V} \right) = n(c_v \ln \Theta + R \ln V) + S_0(n) = \sigma(\Theta, V) + S_0(n) \quad (10)$$

Let us 'divide' the equilibrated system \mathcal{A} in a volume V and at a temperature Θ , or, better said, the whole volume V (or, its whole state space) occupiable, and just occupied now by all its constituents (particles, matter units), with diaphragms (thin infinitely, or, 'thought' only), not affecting thermodynamic properties of \mathcal{A} supposingly, to m parts $\mathcal{A}_i, i \in \{1, m\}, m \geq 1$ with volumes V_i with matter units n_i . Evidently $n = \sum_{i=1}^m n_i$ and $V = \sum_{i=1}^m V_i$.

Let now $S_0(n)=0$ and $S_{0i}(n_i)=0$ for all i . For the entropies S_i of \mathcal{A}_i considered individually, and for the change ΔS , when volumes V, V_i are expressed from the state equations, and for $p=p_i$, $\Theta=\Theta_i$ it will be gained that $\sigma_{[i]}=Rn_{[i]}\ln n_{[i]}$. Then, for $S_i=\sigma_i=n_i(c_v\ln\Theta + R\ln V_i)$ is valid, we have that

$$\begin{aligned} \sum_{i=1}^m S_i &= \sum_{i=1}^m \sigma_i = nc_v \ln \Theta + R \ln \left(\prod_{i=1}^m V_i^{n_i} \right), \\ \Delta S = S - \sum_{i=1}^m S_i &= \sigma - \sum_{i=1}^m \sigma_i = \Delta \sigma = R \ln \frac{V^n}{\prod_{i=1}^m V_i^{n_i}} = -nR \sum_{i=1}^m \frac{n_i}{n} \ln \frac{n_i}{n} > 0 \end{aligned} \quad (11)$$

Let us denote the last sum as B further on, $B < 0$. The quantity $-B$ expressed in (11) is information entropy of a source of messages with an alphabet $[n_1, n_2, n_m]$ and probability distribution $[n_i/n]_{i=1}^m$. Such a division of the system to m parts defines an information source with the information entropy with its maximum $\ln m$.

The result (11), $\Delta S = -nRB$, is a *paradox*, a contradiction with our presumption of not influencing a thermodynamic state of \mathcal{A} by diaphragms, and, leads to that result that the heat entropy S (of a system in equilibrium) is *not* an extensive quantity. But, by the definition of the differential dS , this is *not* true.

Due to this contradiction we must consider a non-zero integrating constants $S_0(n)$, $S_{0i}(n_i)$, in such a way, that the equation $\Delta S = (\sigma + S_0) - \sum_{i=1}^m (\sigma_i + S_{0i}) = 0$ is solvable for the system \mathcal{A} and all its parts \mathcal{A}_i by solutions $S_{0[i]}(n_{[i]}) = -n_{[i]} R \ln(n_{[i]}/\gamma_{[i]})$.

Then $S_{[i]} \triangleq S_{[i]}^{Claus}$, and we write and derive that

$$S^{Claus} = \sum_{i=1}^m S_i^{Claus} = \sum_{i=1}^m n_i R \ln \gamma_i = nR \ln \gamma \Rightarrow \gamma = \gamma_i; \Delta S = 0. \quad (12)$$

Now let us observe an equilibrium, $S^* = S^{Claus} = S^{Boltz} = -kN B^* = -kN \ln N$.

Let, in compliance with the *solution* of Gibbs Paradox, the integration constant S_0 be the (change of) entropy ΔS which is added to the entropy σ to figure out the measured entropy S^{Claus} of the equilibrium state of the system \mathcal{A} (the final state of Gay-Lussac experiment) at a temperature Θ . We have shown that without such correction, the less entropy σ is evidenced, $\sigma = S^{Claus} - \Delta S, \Delta S = S_0$.

Following the previous definitions and results we have

$$\begin{aligned} \Delta S &= \frac{\Delta Q_0}{\Theta} = -nR \ln \frac{n}{\gamma}, \\ \ln \gamma &= \frac{\Delta S}{knN_A} + \ln n = \frac{\Delta S}{kN} + \ln N - \ln N_A, \quad \gamma = N \Rightarrow \frac{\Delta S}{kN} = \ln N_A. \end{aligned} \quad (13)$$

By the entropy ΔS the 'lost' heat ΔQ_0 (at the temperature Θ) is defined.

Thus, our observation can be understood as an information transfer \mathcal{T} in an information channel \mathcal{K} with entropies $H(X)$, $H(Y)$, $H(X|Y)$ and $H(Y|X)$ in (4) but now bound physically; we have these information entropies per one particle of the observed system \mathcal{A} :

$$\begin{aligned}
 \text{input } H(X) &= \frac{\overset{\text{Def}}{S^*}}{kN} = \ln \gamma = -B^* = \ln N = -rB(r) \\
 \text{output } H(Y) &= \frac{\overset{\text{Def}}{\sigma}}{kN} \triangleq -B^{\text{Gibbs}} = -B^{\text{Boltz}} = -B(r), \\
 \text{loss } H(X|Y) &= \frac{\overset{\text{Def}}{S_0}}{kN}, \\
 \text{noise } H(Y|X) &= 0 \text{ for the simplicity;} \\
 H(X|Y) &= -rB(r) - [-B(r)] = -B(r) \cdot (r-1) = (-B^*) \cdot \frac{r-1}{r}, r \geq 1; \frac{1}{r} = \eta_{\max}.
 \end{aligned} \tag{14}$$

For a number m of cells of our railings in the volume V with \mathcal{A} , $m \leq N$ or for the accuracy r of this description of the 'inner structure' of \mathcal{A} (a thought structure of V with \mathcal{A}) and for the number q of diaphragms creating our railings of cells and constructed in such a way that $q \in 1, m-1 >$, we have that $r = (N-1)/q$.

Our observation of the equilibrium system \mathcal{A} , including the *mathematical correction* for Gibbs Paradox, is then describable by the Shannon transfer scheme $[X, \mathcal{K}, Y]$ where

$$H(X) = \frac{S^{\text{Claus}}}{kN}, \quad H(X|Y) = \frac{S_0}{kN}, \quad H(Y) = \frac{S^{\text{Claus}}}{kN}, \quad H(Y|X) = \frac{\Delta S}{kN}. \tag{15}$$

(However, a real observation process described in (15), equivalent to that one with $r=1$, is impossible [6].)

We conclude by that, the diminishing of the measured entropy value about ΔS against S^* awaited, evidenced by **Gibbs Paradox, does not originate in a watched system itself**. Understood this way, **it is a contradiction of a gnozeologic character based on not respecting real properties of any observation** [6]. And, this means the following.

The minimal accuracy of our description of the watched system \mathcal{A} should be for $r = \infty$. In this case we should place $q=0$ diaphragms, no railings is laid, $m=1$, $q=0$. 'We think nothing' about the 'inner structure' of the system \mathcal{A} , for in this case we are not outside of it (for the 'structure' measured - considered is 'created' by 0 diaphragms 'laid down' just from the outside). Thus we define an output information source Y , which is bound physically, and for which its (bound) information entropy is $H(Y) = -B^{\text{Gibbs}} = 0$. Then, the result of such 'observation' is 0, and the loss information entropy is

$$H(X|Y) = \frac{S^*}{kN} = \ln N = H(X).$$

If we consider $m > 1$ as the number of 'windows' being laid down over the measured (observed) system \mathcal{A} from the outside, but now $m = 1$ and then $q = 0$, we resign the possibility of saying about the system \mathcal{A} anything else than that it exists. The whole system \mathcal{A} is now 'viewed' in the just one 'window' but now created by the volume V , the system \mathcal{A} occupies, itself. This only one observation 'window' is 'pressed' to us by the system \mathcal{A} itself and, by this way, 'we ourselves are identical' with it, 'we ourselves are' the system \mathcal{A} . We 'can move' within the volume V , but 'undivided' this time, and thus our measuring could be 'organized' now by the mediation of this only one window, but not laid down from the outside (by us). We are inside of this system (of the volume V) and our measuring is organized with the parameter $r = \infty$. We say that we are identical with the system for we have now no 'step-aside' from it [$H(Y) = H(X) - H(X|Y) = 0$]. This is the reason why we do not see it (from the outside) and, what we can think about it is nothing. In such a sense that we can't 'lay down' the diaphragms over the system ($q \neq 0$ is needed) and create on it our 'windows' ($m > 1$). Because we do not rule our measuring of the system \mathcal{A} we do not 'divide' it by our, just from the outside laid, diaphragms (now $q = 0$) and, for this, we are not able to organize its measuring with the parameter $r < \infty$. 'We ourselves are' the system \mathcal{A} and for this we can't, from the outside, see us, which means we can't see the system. (Or we can evidence its or our certain existence in this window.) The Auto-Reference is not possible; the measuring of the system \mathcal{A} (from the outside !) is intended to be organized in the inside (!) of the system \mathcal{A} itself.³

Our measuring also represents an observation of a certain phenomenon with the *degenerated* probability distribution $[n_i/n]_{i=1}^1$. The information amounts $i(X)$ and $H(X)$ of this phenomenon are equal to 0 and, due to this, our measuring organized by the measuring method with the parameter $m = 1$ or $r = \infty$ respectively will end with the result $H(Y) = 0$ [$H(Y) = \eta_{max} \cdot H(X) - H(X) = 0, \eta_{max} = 1/\infty$].

For this all we need a certain 'step-aside' from the system \mathcal{A} expressed by the value $r < \infty$, but, nevertheless with respecting properties of this 'step-aside', we do not fall into Gibbs Paradox [even when our railings of diaphragms in the mode given $r \in (1, \infty)$ is laid down].

The ideal 'step-aside' is expressed by the value $r = 1$.

The maximal accuracy of our 'description', the accuracy of our watching of the system \mathcal{A} , is achieved for $r = 1$. Then $B(r) = B^*$ and for the output, the input and the loss information entropies it is valid that

$$H(Y) = H(X) = -B^*, \quad H(X|Y) = 0$$

Then we have the ideal 'step-aside' from the system \mathcal{A} . 'We have insight into its inside' and we can draw a layout of measuring its state without the distortion being given by the (at least

³ Our eye cannot see itself by itself only and, even more, it cannot see inside itself by itself only.

partial) transfer (energy) of its state into the organization of our measuring [in the form of the nonzero value of the information loss $H(X | Y)$ when $r \in (1, \infty)$.]

This 'step-aside' enables us to set our measuring with the parameter $q = N - 1, r = 1$. Now we know exactly what we measure, we know that we measure the status of the equilibrium thermodynamic system \mathcal{A} and, by our 'step-aside' from it, we are able to check the precision (the organization) of the measuring method. Just this enables us to organize the measuring of the status of the system \mathcal{A} without nonzero information loss $H(X | Y)$ in the other case being generated by the method itself (for $r \neq 1$), see the mentioned above gnozeological defect.

[But note that something quite different is the realization of the measuring method where the information losses are inevitable yet as the result of the validity of the II. Principle of Thermodynamics (and its equivalents [8]).]

3. Information Thermodynamic Concept Removing Auto-Reference

In the Auto-Reference case, the whole combined machine $\mathcal{O}\mathcal{O}''$ is a system in the equilibrium status. For this status we can introduce the term (quasi)stationary status in which the (infinitesimal) **part of heat is circulating**. Any round of this circulation is lasting the time interval Δt ; infinite, $\Delta t \rightarrow \infty$, for **not ideal model**, or, finite, $\Delta t < \infty$, when the ideal model is used; then the part of heat cannot be the infinitesimal. With the exception of the II. Perpetuum Mobile functionality of this combined machine, which is not possible, see (5) and (6), only the opening of the system and an external activity, a **certain 'step-aside'** between the cycles \mathcal{O} and \mathcal{O}'' , moves it away (prevent it) from this status.

Nevertheless, we sustain our wants of gaining the information (about) $H(X | Y)$ about the structure of the observed \mathcal{O} (the transfer channel \mathcal{K}), we want the nonzero value ΔA^* , the nonzero information $H^*(Y^*) = \Delta A^* / kT_W > 0$.

Then, necessarily, the mirror, the reverse Carnot Cycle \mathcal{O}'' (the transfer channel \mathcal{K}'') is to be constructed with that 'step-aside' (excluding that stationarity) from the observed $\mathcal{O} \cong \mathcal{K}$. Now we mean that the 'step-aside' of the *observing* process \mathcal{O} from the *observed* process \mathcal{O}'' is realized by the difference $T_W - T''_w > 0$. Now, within this thermodynamic point of view, it is valid that $\Delta A'' < \Delta A^*$ for $T_0 = T''_w, T''_w \triangleq T_0^*$

$$\Delta A'' = \Delta Q''_W \cdot \left(1 - \frac{T_0}{T''_W}\right) = \Delta Q_W \cdot \frac{T''_W}{T_W} \cdot \left(1 - \frac{T_0}{T''_W}\right) \quad (16)$$

Then, for the whole information amount $\Delta A^* / kT_W$ of our combined cycle it is valid that

$$\frac{\Delta A^*}{kT_W} = H(Y'') - H(Y'') \cdot \beta^* = H(Y'') \cdot \left(1 - \frac{T_0}{T_W}\right) = H(Y'') \cdot \left[1 - \frac{H(X|Y)}{H(X''|Y'')}\right] \quad (17)$$

The structure $H(X | Y)$ of the observed transfer (channel, process) $\mathcal{O} \cong \mathcal{K}$ is measurable with the 'step-aside' only, created now by different temperatures ($T_W > T''_w$). The result is

$$\frac{\Delta A^*}{kT_W} > 0, \left[\frac{\Delta A^*}{kT_W} \cong f[H(X | Y)] > 0 \right] \quad (18)$$

Following (5), (6) and (9) the Auto-Reference arises just when

$$T_W = T''_w \left[\Rightarrow H(Y) = 0 \right] \quad (19)$$

Now we will describe, in the *information thermodynamic way*, the problem of a Self-Observation or the Auto-Reference problem and will draw a method of its ceasing.

For we want to have the information $H(X | Y)$ describing the structure of the transfer process $\mathcal{O} \cong \mathcal{T}$ which we observe we need gain a nonzero value (difference) ΔA^* and, consequently, a nonzero information $H^*(Y^*)$,

$$H^*(Y^*) = \frac{\Delta A^*}{kT_W} > 0 \quad (20)$$

Then we need a 'mirror', the reverse Carnot Cycle $\mathcal{O}'' \cong \mathcal{T}''$ (or the relevant transfer channel \mathcal{K}'') would be constructed in such a way that the mentioned 'step-aside' from the observed transfer channel \mathcal{K} was respected. [It is the 'step-aside' of the observing process (\mathcal{O}'' , \mathcal{T}'') from the observed process (\mathcal{O} , \mathcal{T}); also we can consider a computing process $\bar{\kappa}$ and its description - the program $\bar{\eta}$, and its observation, see later].

Now the required 'step-aside' is realized by the temperature difference $T_W - T''_w > 0$. Thus now we consider, within the frame of our thermodynamic approach, that $\Delta A'' < \Delta A^*$ for $T_0 = T''_0$ and then, under the condition $\Delta Q_0 = \Delta Q''_0$, it will be valid for the cycles \mathcal{O} , \mathcal{O}'' and $\mathcal{O}\mathcal{O}''$ that

$$\begin{aligned} \Delta Q''_W &= \Delta A'' + \Delta Q''_0 \\ &= \Delta Q''_W \cdot \eta''_{max} + \Delta Q''_W \cdot (1 - \eta''_{max}) \\ &= \Delta Q''_W \cdot \eta''_{max} + \Delta Q''_W - \Delta Q''_W \cdot \eta''_{max} \end{aligned} \quad (21)$$

but also it is valid that

$$\begin{aligned} \Delta Q_W &= \Delta Q''_W \cdot \eta''_{max} + \Delta Q''_W \cdot (1 - \eta''_{max}) \\ &= \Delta Q''_W \cdot \eta''_{max} + \Delta Q''_W - \Delta Q''_W \cdot \eta''_{max} \end{aligned} \quad (22)$$

From the proposition $\Delta Q_0 = \Delta Q''_0$ [from relations (21) and (22)] pro $\Delta Q''_W$ follows that

$$\Delta Q_W \cdot (1 - \eta_{max}) = \Delta Q_W'' \cdot (1 - \eta_{max}'') \quad (23)$$

With the denotation $\beta \triangleq (1 - \eta_{max})$, $\beta'' \triangleq (1 - \eta_{max}'')$ we write

$$\begin{aligned} \Delta Q_W \cdot \beta &= \Delta Q_W'' \cdot \beta'' \\ \frac{\Delta Q_W}{\Delta Q_W''} &= \frac{T_W}{T_W''} = \frac{\frac{T_0}{T_W}}{\frac{T_0}{T_W''}} = \frac{\beta''}{\beta} \\ &\Rightarrow \\ \Delta Q_W'' &= \Delta Q_W \cdot \frac{\beta}{\beta''} = \Delta Q_W \cdot \frac{T_W''}{T_W}, \Delta Q_W'' < \Delta Q_W \end{aligned} \quad (24)$$

Within the cycles \mathcal{O} and \mathcal{O}'' the following relations are valid,

$$\begin{aligned} \frac{T_W - T_0}{T_W} &> \frac{T_W'' - T_0}{T_W''}, T_W > T_W'' \\ \frac{T_0}{T_W''} > \frac{T_0}{T_W} \text{ a tedy } \frac{Q_0}{Q_W''} &> \frac{Q_0}{Q_W}, Q_W > Q_W'' \\ &\Rightarrow \\ \frac{H(Y''|X'')}{H(Y'')} &> \frac{H(X|Y)}{H(X)} \end{aligned} \quad (25)$$

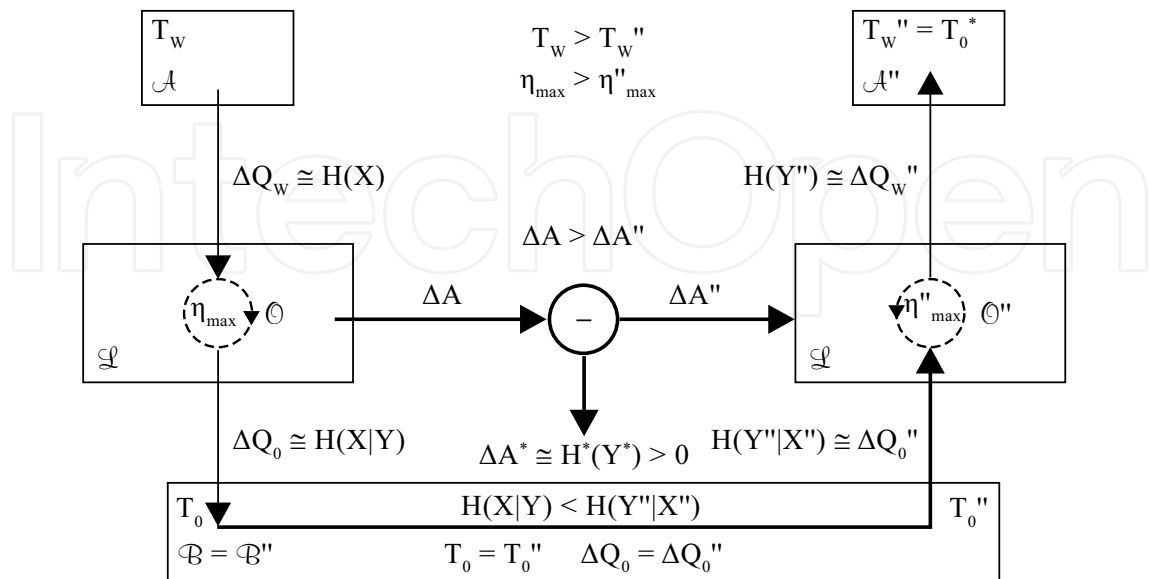


Figure 2. The concept for ceasing the Auto-Reference

By (23) and (24) for $\Delta A''$ it is valid in the cycle \mathcal{O}'' that

$$\begin{aligned}\Delta A'' &= \Delta Q_W'' \cdot \left(1 - \frac{T_0}{T_W''}\right) = \Delta Q_W \cdot \frac{T_W''}{T_W} \cdot \left(1 - \frac{T_0}{T_W''}\right) = \\ &= \Delta Q_W \cdot \left(\frac{T_W''}{T_W} - \frac{T_0}{T_W}\right) = k \cdot H(X) \cdot (T_W'' - T_0) \\ &= k \cdot H(X) \cdot T_W'' \left(1 - \frac{T_0}{T_W''}\right) = k \cdot H(X) \cdot T_W'' (1 - \beta'') = k \cdot T_W'' \cdot H(Y'')\end{aligned}\quad (26)$$

and, further, for ΔA in the cycle \mathcal{O} we have

$$\Delta A = k \cdot H(X) \cdot T_W (1 - \beta) = k \cdot H(X) \cdot T_W \left(1 - \frac{T_0}{T_W}\right) \quad (27)$$

and thus, for the cycles \mathcal{O}'' and \mathcal{O} it is valid that

$$\frac{\Delta A''}{k T_W''} = H(X) \cdot \left(1 - \frac{T_0}{T_W''}\right) = H(X) \cdot (1 - \beta'') = H(X) \cdot \eta_{max}'' \quad (28)$$

$$\frac{\Delta A}{k T_W} = H(X) \cdot \left(1 - \frac{T_0}{T_W}\right) = H(X) \cdot (1 - \beta) = H(X) \cdot \eta_{max}$$

For the whole work ΔA^* of the combined cycle $\mathcal{O}\mathcal{O}''$ we have

$$\Delta A^* = \Delta A - \Delta A'' = \left[k T_W \cdot H(X) \cdot (1 - \beta) - k T_W'' \cdot H(X) \cdot (1 - \beta'') \right] > 0 \quad (29)$$

Then, for the *whole change* of the thermodynamic entropy within the combined cycle $\mathcal{O}\mathcal{O}''$ (measured in information units *Hartley, nat, bit*) and thus for the change of the whole information entropy $H^*(Y^*)$ [$\triangleq H_C$], following the relation (29), it is valid

$$\begin{aligned}H^*(Y^*) &= \frac{\Delta A^*}{k T_W} = H(X) \cdot \left[(1 - \beta) - \frac{T_W''}{T_W} \cdot (1 - \beta'') \right] \\ &= H(X) \cdot \left(1 - \frac{T_0}{T_W} - \frac{T_W''}{T_W} + \frac{T_0}{T_W}\right) = H(X) \cdot \left(1 - \frac{T_W''}{T_W}\right)\end{aligned}\quad (30)$$

It is valid, for ΔA^* is a *residuum work* after the work ΔA has been performed at the temperature T_w . Evidently, the sense of the symbol T''_w (within the double cycle $\mathcal{O}\mathcal{O}''$ and when $\Delta Q_0 = \Delta Q''_0$) is expressible by the symbol T_0^* , which is possible, for the working temperatures of the whole cycle $\mathcal{O}\mathcal{O}''$ are T_w and $T''_w = T_0^*$. The relation (30) expresses that fact that the double cycle $\mathcal{O}\mathcal{O}''$ is the direct Carnot Cycle just with its working temperatures $T_w > T''_w = T_0^*$. In the double cycle $\mathcal{O}\mathcal{O}''$ it is valid that

$$\beta'' = \frac{\Delta Q_0''}{\Delta Q_w''} = \frac{\frac{\Delta Q_0''}{T_w''}}{\frac{\Delta Q_w''}{T_w''}} = \frac{H(Y''|X'')}{H(Y'')} = \frac{T_0}{T_w''}, T_w'' = T_0^*, \text{ cyklus } \mathcal{O}''$$

$$\beta = \frac{\Delta Q_0}{\Delta Q_w} = \frac{\frac{\Delta Q_0}{T_w}}{\frac{\Delta Q_w}{T_w}} = \frac{H(X|Y)}{H(X)} = \frac{T_0}{T_w}, \text{ cyklus } \mathcal{O} \tag{31}$$

$$\frac{\beta}{\beta''} = \frac{T_w''}{T_w} = \frac{T_0^*}{T_w} \triangleq \beta^*$$

and then, by (30) a (31) is writable that

$$\frac{\Delta A^*}{kT_w} = H(X) \cdot (1 - \beta^*) = H(X) \cdot \left[1 - \frac{H(X|Y) \cdot H(Y'')}{H(Y''|X'') \cdot H(X)} \right] > 0 \tag{32}$$

It is ensured by the propositions $T_w > T''_w$, $T''_0 = T_0$ and also by that fact, that the loss entropy $H(X|Y)$ is described and given by the heat $\Delta Q_0 = \Delta Q''_0$. But by our combined cycle $\mathcal{O}\mathcal{O}''$ it is valid too that

$$H(X) = \frac{\Delta Q_w}{kT_w} = \frac{\Delta Q_w''}{kT_w''} = H(Y'') \left[= \frac{\Delta Q_w''}{kT_0^*} \right] \tag{33}$$

and by both parts of (4) we have

$$\frac{H(X|Y)}{H(Y''|X'')} = \beta^* < 1 \tag{34}$$

For the whole information entropy $\Delta A^*/kT_w$ (the whole thermodynamic entropy \mathcal{S}_C in information units) and by following the previous relations also it is valid that

$$\begin{aligned} \frac{\Delta A^*}{kT_w} &= H(Y'') - H(Y'') \cdot \beta^* = H(Y'') \cdot \left(1 - \frac{T_0^*}{T_w}\right) \\ &= H(Y'') \cdot \left[1 - \frac{H(X|Y)}{H(X''|Y'')}\right] \end{aligned} \quad (35)$$

And thus, the structure of the information transfer channel \mathcal{K} (expressed by the quantity $H(X|Y)$) is measurable by the value $H^*(Y^*)$ from (20), (32) and (35). Symbolically, we can write, using a growing function f ,

$$H^*(Y^*) = \frac{\Delta A^*}{kT_w} \cong f[H(X|Y)] > 0 \quad (36)$$

The cycles \mathcal{O} , \mathcal{O}'' and $\mathcal{O}\mathcal{O}''$ are the Carnot Cycles and thus, from their definition and construction they are, imaginatively⁴ in principle, the **infinite** cycles; in each of them the following *criterion of an infinite cycle* (see [14]) it is valid inevitably,

$$T(X^{[1]}; Y^{[1]}) = H(X^{[1]}) - H(X^{[1]}|Y^{[1]}) = H(Y^{[1]}) > 0 \quad \text{and} \quad \Delta S_c^{[1]} = 0 \quad (37)$$

The construction of the cycle $\mathcal{O}\mathcal{O}''$ enables us to recognize that the *infinite cycle* \mathcal{O} is running; the unsolvable case of the Auto-Reference in $\mathcal{O}\mathcal{O}''$ occurs just when and only when

$$T_w = T_w'', \quad \left[[T_w = T_w''] \Rightarrow [H^*(Y^*) = 0] \right] \quad (38)$$

(Thus, when we contemporarily await anything else than the zero output or an output which is not relevant to the given input, by this only awaiting, we require a construction of the Perpetuum Mobile functionality.)

Our observation of the process \mathcal{O} now being considered as a model or as the realization of the computing process $\bar{\kappa}$ in a certain Turing Machine TM or an information transfer \mathcal{T} in a certain channel \mathcal{K} - the measuring of a transferring structure expressed by the value of the quantity $H(X|Y)$, is thus possible but only by another process with a certain 'step-aside' from the observed process \mathcal{O} in-built, with a certain 'mirror' $\mathcal{O}'' \cong \bar{\mathcal{T}}$, with the aids of another transferring structure expressed by the value of the quantity $H(X''|Y'')$.

From the both processes, the cycles \mathcal{O} and \mathcal{O}'' , the whole combined and one cycle $\mathcal{O}\mathcal{O}''$ is created to be modeling the whole and one transfer channel $\mathcal{K}\mathcal{K}''$ in which the observed, the measured property, the value $H(X|Y)$ in the given 'time click' p (the click, the interval) expressing the

⁴ When an infinite reserve of energy would exist.

structure of the channel \mathcal{K} is one of the ingoing (input) messages (informations). The resulting double cycle $\mathcal{O}\mathcal{O}''$ performs as a direct Carnot Cycle with the working temperatures T_W and $T_{0'}^*$, $T_W > T_{0'}^*$, $T_w'' = T_0^*$.

The required 'step-aside' is created by the difference $T_W - T_0^* > 0$. The whole information entropy (the thermodynamic entropy in information units) of the whole isolated system in which our double cycle $\mathcal{O}\mathcal{O}''$ is running, the whole output information $H^*(Y^*)$ of this double cycle $\mathcal{O}\mathcal{O}''$ is then a certain function $f(\cdot)$ of the *measure* $H(X \mid Y)$ of the structure being measured (observed) and as such, of the value of the argument $H(X \mid Y)$ from the relation (36).

Remark: Also the following consideration is possible: We use the change $H(Y)$ of the output entropy of the observed process \mathcal{O} as its reaction to a change $H(X)$ of the input entropy and just by the evidenced $H(X \mid Y)$. Our measuring is then characterized, in the same sense as in (36) and (37), by the whole value

$$H(X) \cdot (1 - \beta) - H(X) \cdot (1 - \beta'') = H(X) \cdot \frac{T_0}{T_w''} \cdot \left(1 - \frac{T_w''}{T_W} \right) > 0$$

Now it is obvious that the transferring (copying, measuring, observing) of the structure of the channel \mathcal{K} is possible then and only then a certain structural 'step-aside' between the observed object (the transferred or the measured structure, now and here the structure of a transfer channel \mathcal{K}) and the observing process (the channel \mathcal{K}'' with the transferring process \mathcal{T}'') expressed by the nonzero and positive values of the difference (2) is possible.

By the term 'step-aside' of the **observing** computing process (let us denote it as $\vec{\kappa}''$) from the **observed** computing process (let us denote it as $\vec{\kappa}$) we understand a **time delay** between them, better said, it is a **tracing**⁵ of the observed computing process $\vec{\kappa}$.

4. Turing Computing, Auto-Reference and Halting Problem

Turing Machine (TM) is driven by a *program* which is interpreted by its *Control Unit (CU_{TM})*. The Control Unit CU_{TM} is a *finite automaton* (Mealy's or Moore's *sequential machine*). The *program* for the TM consists of the finite sequence $\vec{\eta}$ of instructions $\eta_{[.]}$

$$\vec{\eta} = \left(\eta_q \right)_{q=1}^{q \in \mathbb{N}} = \left[\left(s_i, x_k, s_j, y_l, D \right) \right]_{q=1}^{q \in \mathbb{N}}, \quad |\vec{\eta}| \in \mathbb{N} \quad (39)$$

Each of these instructions describes an overwriting rule of a *regular grammar* [15, 19, 21],

⁵ In the programmers' manner or language: listing, cross-reference, memory dump.

$$s_i \rightarrow (x_k, y_l, D) s_j \quad (40)$$

performed in the given step (time, moment) p , $p \in \mathbb{N}$, of the TM 's activity;

- s_i is the i -th *nonterminal symbol* of the regular grammar, or, respectively, it is a *status* of the CU_{TM} within the actual step $p \in \mathbb{N}$ of the TM 's activity,
- x_k is an *input terminal symbol* being read from the *input-output tape* of the TM within the actual step p of the TM 's activity,
- y_l is an *output terminal symbol* by which the CU_{TM} overwrites the symbol x_k which has been read (in the actual step p of the TM 's activity),
- s_j is the *successive status* of the CU_{TM} , given by the instruction for the following step $p + 1$ of the CU_{TM} .

Within the actual step p of the TM 's activity the CU_{TM} is changing its status to s_j [this change is based on the status s_i , and the symbol x_k has been read ($s_i \xrightarrow{x_k} s_j$)], and is performing the transformation

$$x_k \rightarrow y_l \quad (41)$$

on the scanned (actual) *position* of the input-output tape,

- D determines the *moving direction* of the *read-write head* of the CU_{TM} after the symbol y_l has been recorded [in the status s_{j_p} (s_{j_p} denotes s_j for the step p) used further on as the following one, $s_{j_p} \stackrel{\text{Def}}{=} s_{i_{p+1}}$], $D \in \{L, R\}$.

The value L or R of the symbol D determines the *left slip* or the *right slip* from the actual position on the input-output tape to its (left or right) neighbor after the transformation x_k to y_l has been performed.

The *oriented edge* of the *transition graph* of the CU_{TM} (the finite automaton) is described by the symbol $s_i \xrightarrow{(x_k, y_l, D)} s_j$. The TM 's activity generates a sequence of the instructions *having been performed* in steps p , $[(s_{i_p}, x_{k_p}, s_{j_p}, y_{l_p}, D_p)]_{p=1}^{p=p_{\text{last}}}$,

$$s_{i_p} \xrightarrow{(x_{k_p}, y_{l_p}, D_p)} s_{j_p}, \text{ further on } s_{j_p} = s_{i_{p+1}} \quad (42)$$

(the edge of the oriented transition graph of the CU_{TM} in the step p), by which the *computing process* ($\vec{\kappa}$) has gone through (from the first step $p=1$ till, *for this while*, the last step $p=p_{\text{last}}$ of the

TM 's activity). They are also the *overwriting rules* of the regular grammar, being performed within each step $p, p \geq 1$,

$$s_{i_p} \rightarrow \left(x_{k_p}, y_{l_p}, D_p \right) s_{j_p}, s_{i_p} = s_{i_{p+1}} \quad (43)$$

By this way a *regular language* of the words (x_{k_p}, y_{l_p}, D_p) or, respectively, a regular language of the instructions $(s_{i_p}, x_{k_p}, s_{j_p}, y_{l_p}, D_p)$ having been performed is defined. This second regular language is describable by the rules (of a regular grammar)

$$S_{i_p} \rightarrow \left(s_{i_p}, x_{k_p}, s_{j_p}, y_{l_p}, D_p \right) S_{j_p}, S_{i_p} \stackrel{\text{Def}}{=} S_{i_{p+1}} \quad (44)$$

being applied in each step $p \geq 1$ of the TM 's activity. Thus, this language is to be acceptable by a certain finite automaton with n states $S_{[.]}$.

When this language is *infinite*⁶(the infinite chain of instructions of the finite length), such its word

$$\left[\left(s_{i_1}, x_{k_1}, s_{j_1}, y_{l_1}, D_1 \right), \dots, \left(s_{i_p}, x_{k_p}, s_{j_p}, y_{l_p}, D_p \right) \right]_{p=l} \quad (45)$$

of the length l exists that for that finite automaton [with n states $S_{[.]}$ and the transition rules (42) or (44)] the **Pumping Lemma** [19, 21] **it is valid**

$$n \leq l < 2n \quad (46)$$

4.1. Auto-Reference in Turing Computing

Although any instruction of the Turing Machine TM describes one step of the *computing process* in this TM, it is considerable as a *description (of one step) of an information transfer process* running in a certain transfer channel \mathcal{K} ; we postulate the relation $TM \cong \mathcal{K}$. The computing process in the TM is, also, a transfer process in a channel \mathcal{K} . For $\mathcal{K} \cong \mathcal{O}$ it is valid that $TM \cong \mathcal{O}$.

In each step $p > 1$ of its activity, the $TM \cong \mathcal{K}$ is accepting *its own configuration from the previous step* $p - 1$ as its input, includes its contemporary status $(s_{i_p} = s_{j_{p-1}})$ and generates its status $[s_{i_{(p+1)}}]$ and the configuration for the next step $p + 1$, etc.⁷

⁶ Better said, having the *arbitrary* (but finite) length.

⁷ $(s_{i_p}, x_{k_p}, s_{j_p}, y_{l_p}, D_p)(\vec{\sigma}_p, s_{p'}, \vec{\rho}_p) \rightarrow (\vec{\sigma}_{p+1}, s_{i_{p+1}}, \vec{\rho}_{p+1})$, see further.

Similarly it is valid for the configurations (denoted now by X_p and Y_p), see further on.

For each $p \geq 1$ we consider the actual instances of the stochastic quantities⁸ X, Y ,

$$\begin{aligned} X \triangleq X_p, Y \triangleq Y_p; X|Y \triangleq X_p|Y_p, Y|X \triangleq Y_p|X_p; Y_p = X_{p+1} \\ X_p|X_{p+1} \triangleq X_p \hat{\circ} X_p X_{p+1}, X_{p+1}|X_p \stackrel{\Delta}{=} X_{p+1} \hat{\circ} (X_p X_{p+1})^{-1} \end{aligned} \tag{47}$$

In any step p of the TM 's activity its own configurations $(\vec{\sigma}_p, s_p, \vec{\rho}_p)$ - members of the sequence - of the **computing process** $\vec{\kappa} \stackrel{\text{Def}}{=} [(\vec{\sigma}_p, s_p, \vec{\rho}_p)]_{p=1}^p$, can be considered as follows;

- let now the stochastic quantity X_p be realized by the chain

$$(\vec{\sigma}_p, s_p, \vec{\rho}_p) \in \mathbf{T}^* \times \mathbf{S} \times \mathbf{T}^*; p = 1, \vec{\sigma}_1 = \vec{\varepsilon} \text{ a } \vec{\rho}_1 = \vec{\xi}, s_p = s_0 \tag{48}$$

- let now the stochastic quantity Y_p be realized by the chain

$$[\vec{\sigma}_{p+1}, s_{p+1}, \vec{\rho}_{p+1}] \in \mathbf{T}^* \times \mathbf{S} \times \mathbf{T}^* \tag{49}$$

Then, the computing process in the $TM \cong \mathcal{K}$ is describable informationally,^{9,10}

$$\begin{aligned} H(X) \triangleq H(X_p) &= H(\vec{\sigma}_p, s_p, \vec{\rho}_p) \\ H(Y) \triangleq H(Y_p) &= H[X_{p+1}] = H[\vec{\sigma}_{p+1}, s_{p+1}, \vec{\rho}_{p+1}] \\ H(X|Y) \triangleq H(X_p|Y_p) &= H[X_p|X_{p+1}], H(Y|X) \triangleq H(Y_p|X_p) = H[X_{p+1}|X_p] \\ T(X;Y) \triangleq T(X_p;Y_p) &= H(X_p) - H[X_p|X_{p+1}] \\ &= H(\vec{\sigma}_p, s_p, \vec{\rho}_p) - H[(\vec{\sigma}_p, s_p, \vec{\rho}_p) | [\vec{\sigma}_{p+1}, s_{p+1}, \vec{\rho}_{p+1}]] \\ T(Y;X) \triangleq T(Y_p;X_p) &= H(X_{p+1}) - H[X_{p+1}|X_p] \\ &= H[\vec{\sigma}_{p+1}, s_{p+1}, \vec{\rho}_{p+1}] - H[\vec{\sigma}_{p+1}, s_{p+1}, \vec{\rho}_{p+1} | (\vec{\sigma}_p, s_p, \vec{\rho}_p)] \end{aligned} \tag{50}$$

The Auto-Reference arises with the following description of the computing (transfer, observation) process when, e.g., for a certain $p \geq p^* \geq 1$,

⁸ '≡' now denotes the *substring* from the beginning of the string.

⁹ The transitions are given by the η_{q_p} called by $X_p X_p | X_{p+1} \rightarrow \eta_{q_p} [\eta_{q_p}^{-1}(X_{p+1}) = X_p]$.

¹⁰ $\eta_{q_p}^{-1}(X_{p+1}) = X_p$ - comparison of the structures of the X_p and the X_{p+1} (in $X_p | X_{p+1}$).

$$\begin{aligned}
 H(X_p) - H[X_p | X_{p+1}] &= H(X_{p+1}) - H[X_{p+1} | X_p], \quad p \geq 1 \\
 H(X_p) &= H(X_p | X_{p+1}); \quad X_p = X_p \hat{\circ} X_{p+1}, \quad X_{p+1} = \varepsilon \\
 H(X_{p+1}) &= H(X_{p+1} | X_p), \quad [H[X_{p+1} | X_p] = 0]
 \end{aligned}
 \tag{51}$$

This way of considerations 'constructs' the TM 's infinite cycle from the programmer's point of view, Self-Observation in an information point of view and a stationary status from the thermodynamics point of view.

In any case a 'step-aside' to gain something else than the zero output is required.

By the 'step-aside' of the *observing* computing process from the *observed* computing process we mean a *time delay* between those two processes or, better said, a *staging of the trace* of the observed process.

4.2. Halting Problem as Auto-Reference

Now we are considering a certain TM (the observed machine) being driven by a program $\vec{\eta}$ and working with a certain input word $\vec{\xi}$. Let this activity be described by the word $[d(TM)]$.

Let us consider that the TM with the input word $\vec{\xi}$

- halts, $HALT_{TM}$, whether the word $\vec{\xi}$ is accepted or rejected,

$$HALT_{TM} = \{ HALT_{TM}^{Accept} \cup HALT_{TM}^{Reject} \}
 \tag{52}$$

- does not halt, $LOOP_{TM}^{\infty}$ (the TM 's infinite cycle)

Let us construct the three new Turing Machines M_1 , M_2 and M_3 as follows [18]¹¹

- M_1 works with the input word $[d(TM) * \vec{\xi}]$ in that way that

- halts, $HALT_{M_1}^{Accept}$,

- stops, $HALT_{M_1}^{Reject}$,

$$HALT_{M_1}^{Accept} \Leftarrow HALT_{TM}, \quad HALT_{M_1}^{Reject} \Leftarrow LOOP_{TM}^{\infty}
 \tag{53}$$

- M_2 **modifies** the activity of the M_1 in that way, that the input word which is being worked with is $[d(TM) * \vec{\xi}]$ and

- halts, $HALT_{M_2}'$

¹¹ Minski's proof for the undecidability of the Halting Problem (Entscheidungsproblem type).

- does not halt, $LOOP_{M_2}^\infty$,

$$\begin{aligned} & \text{HALT}_{M_2} \Leftarrow LOOP_{TM}^\infty \left[\Rightarrow \text{HALT}_{M_1}^{\text{Reject}} \right] \\ & LOOP_{M_2}^\infty \Leftarrow \text{HALT}_{TM} \left[\Rightarrow \text{HALT}_{M_1}^{\text{Accept}} \right] \end{aligned} \tag{54}$$

- M_3 is an 'extension' of the M_2 : it doubles its own input word $[d(TM)]$ into $[d(TM) * d(TM)]$ and gives it to the input (of its sub-machine) M_2 and
- halts, HALT_{M_3} ,
- does not halt, $LOOP_{M_3}^\infty$,

$$\begin{aligned} & \text{HALT}_{M_3} \equiv \text{HALT}_{M_2} \Leftarrow LOOP_{TM}^\infty \left[\Rightarrow \text{HALT}_{M_1}^{\text{Reject}} \right] \\ & LOOP_{M_3}^\infty \equiv LOOP_{M_2}^\infty \Leftarrow \text{HALT}_{TM} \left[\Rightarrow \text{HALT}_{M_1}^{\text{Accept}} \right] \end{aligned} \tag{55}$$

- But, when the machine $M_3 \equiv TM$ accepts the description $d(M_3)$, thus it is valid that $d(M_3) \equiv d(TM)$, then

$$\begin{aligned} & \left[\text{HALT}_{TM} \Leftarrow LOOP_{TM}^\infty \right] \wedge \left[LOOP_{TM}^\infty \Leftarrow \text{HALT}_{TM} \right] \\ & \equiv \\ & \text{HALT}_{TM} \Leftrightarrow LOOP_{TM}^\infty \end{aligned} \tag{56}$$

This result (56) is the **contradiction**. It is the consequence of the Cantor diagonal argument having been used carrying the Auto-Reference to the sequence of the machines (TM, M_1, M_2, M_3) , or, respectively, to the sequence of the machines (TM, M_3) ,

$$\left(\underline{TM}, M_3 \equiv \underline{TM} \right) \tag{57}$$

and is leading us to that opinion that the proposition about the decidability of the Halting Problem (recognizing the $LOOP_{TM}^\infty$ state) is not right.

In any given step $p \geq 1$ the machine TM is deciding about itself (it is working with the description of its own actual status), it is the channel 'transferring' its own structure, it is the Self-Observer. Thus it is in a stationary status in the thermodynamic point of view.

*Within this point of view, we can envisage two identical, but reversed mutually, ideal Carnot Cycles connected together. In this sense, these two machine \mathcal{O} and \mathcal{O}'' create the equilibrium system $\mathcal{O}\mathcal{O}''$, in which we introduce the term **stationary status**.*

Within such a system the (infinitesimal) **part of heat is circulating** through the whole combined machine $\mathcal{O}\mathcal{O}''$. Let this fact be now the **thermodynamic model of the infinite cycle** being started by the Self-Observation, by the Auto-Reference action (56), (57); one run is like an uninteruptable operation.

The recursive call of the function $\overrightarrow{d(TM)}$ (of the machine \overrightarrow{TM}) by the same function $\overrightarrow{d(TM)}$ (by the machine TM) with the same argument $[d(TM) * d(TM)]$ is now given. The Auto-Reference (56), (57) is then the generative function for the infinite sequences, nevertheless thought only, as the consequence of the stationarity concept in-built in this type of consideration,

$$\begin{aligned}
 (TM, TM, \dots, TM, \dots) &\triangleq \left(TM_{\Delta t_p} \right)_{p=1}^{\infty} \\
 [d(TM), d(TM), \dots, d(TM), \dots] &\triangleq \left[[d(TM)]_{\Delta t_p} \right]_{p=1}^{\infty} \\
 [HALT_{TM}^{\infty}, HALT_{TM}^{\infty}, \dots, HALT_{TM}^{\infty}, \dots] &\triangleq \left[(HALT_{TM}^{\infty})_{\Delta t_p} \right]_{p=1}^{\infty} \\
 [LOOP_{TM}^{\infty}, LOOP_{TM}^{\infty}, \dots, LOOP_{TM}^{\infty}, \dots] &\triangleq \left[(LOOP_{TM}^{\infty})_{\Delta t_p} \right]_{p=1}^{\infty}
 \end{aligned} \tag{58}$$

We envisage, within this *time-expansion* of the (56) or (57) [which possibility follows from the (quasi)stationarity concept], the infinite cycle in the observed machine TM arises, based on its Self-Description $[\vec{\xi}] \equiv [d(TM)]$. But, following the Auto-Reference construction, it 'runs' in the double-machine $(TM, M_3 \equiv TM) \cong \mathcal{O}\mathcal{O}''$.¹²

The Auto-Reference step that is to solve the Halting solve the Halting Problem proves, only, its own disusability **creates just a certain image of what is to be possibly discovered - the infinite cycle** in the form of the infinite constant time sequences [when the time expansion (58) for $p \geq 1$ is considered].

As it is valid for any **stationary status**, also this one can be ceased or **can be excluded** by an external action, by the 'step-aside', **by the staging** as follows.

5. Concept Removing Halting Problem

We suppose that in the case of a computing process running in a TM its status $LOOP_{TM}^{\infty}$ (the infinite cycle) is decidable within the Observing Turing Machine (OTM) by using the TM 's

¹² Generally, the cycle $\mathcal{O}\mathcal{O}''$ is considered as the reversible only.

trace. By this trace, the machine *OTM* generates and controls the 'combined observing process' for the process in the *TM*.

We will show and use the fact, that certain *regular sequences are generated. If they are infinite, they are, inevitably, periodical*; as such, they are *decidable languages for their infinity* [1].

We will use the alphabet of terminal symbols $T = \{I, B\}$ and these structures:

- (s_i, x_k, s_j, y_l, D) is the *instruction*
- $(\vec{\sigma}, s_{[\cdot]}, \vec{\rho})$ is the *configuration*
- $(\varepsilon[\sigma, s_{[\cdot]}, \rho] \varepsilon)$ is the *configuration type*

Further, we introduce the *general configuration type X* ;

- $X = (\vec{\sigma}, s_{[\cdot]}, \vec{\rho}) \triangleq (\mathbf{B} \sigma, s_{[\cdot]}, \rho \mathbf{B})$.

By this general configuration type the chains, e.g. $\overrightarrow{\mathbf{BIB}}_{s_{[\cdot]}} \overrightarrow{\mathbf{BIB}}$ are ment,

$$\begin{aligned} \overrightarrow{\mathbf{BIB}}_{s_{[\cdot]}} \overrightarrow{\mathbf{BIB}} &= \text{BBBBIIIBBBBBIIIBBB}_{s_{[\cdot]}} \text{BBIIBBBBBBBBB} \\ &\in (\vec{\sigma}, s_{[\cdot]}, \vec{\rho}) \end{aligned}$$

The computing process in the observed TM generates the grammar of a regular language of instructions and, also, of general types of configurations especially, infinite possibly, and thus cyclical. As such, they are decidable languages for their infinity.

These grammars are given by the initial input $\vec{\xi}$, or, respectively, by the initial configuration $(\varepsilon, s_0, \vec{\xi})$ and, also, by the instructions η_{q_p} of the programme $\vec{\eta}$, being generated by the sequence of the steps p of the *TM* 's activity in which the *TM* instructions are interpreted. This sequence itself is pressed out by the configurations having been generated (See Appendix).

The Auto-Reference arises when, e.g., for a certain $p \geq p^* \geq p^0 \geq 1$,

$$\begin{aligned} H(\overrightarrow{\mathbf{X}}_p) - H(\overrightarrow{\mathbf{X}}_p \mid \overrightarrow{\mathbf{X}}_{p+1}) &= H(\overrightarrow{\mathbf{X}}_{p+1}) - H(\overrightarrow{\mathbf{X}}_{p+1} \mid \overrightarrow{\mathbf{X}}_p) = 0 \\ \overrightarrow{\mathbf{X}}_p &\triangleq (\mathbf{X}_{p^*}, \mathbf{X}_{p^*+1}, \dots, \mathbf{X}_p) \\ \overrightarrow{\mathbf{X}}_{p+1} &\triangleq (\overrightarrow{\mathbf{X}}_p, \mathbf{X}_{p+1}, \dots, \mathbf{X}_{2p+2-p^*}) \equiv (\overrightarrow{\mathbf{X}}_p \overrightarrow{\mathbf{X}}_p) \end{aligned} \tag{59}$$

Also we can write [the similar is writable for (22), (23); also see the remark 7].

$$\begin{aligned}
 \overline{X}_p - (\overline{X}_p \hat{\circ} \overline{X}_{p+1}) &= \overline{X}_{p+1} - (\overline{X}_p \hat{\circ} \overline{X}_{p+1}^{-1}) = \varepsilon \\
 \overline{X}_p &= (\overline{X}_p \hat{\circ} \overline{X}_{p+1}) = (\overline{X}_p \hat{\circ} \overline{X}_{p+1}^{-1}) \\
 H(\overline{X}_p) &= (\overline{X}_p \hat{\circ} \overline{X}_{p+1}) = H(\overline{X}_p \hat{\circ} \overline{X}_{p+1}^{-1}) = H(\overline{X}_{p+1} \hat{\circ} \overline{X}_p) = 0 \\
 [H(\overline{X}_{p+1}) = 0] &\Leftrightarrow [Pr(\overline{X}_p) = Pr(\overline{X}_p \overline{X}_p) = Pr([\overline{X}_p]^+) = 1]
 \end{aligned}
 \tag{60}$$

where $Pr(\cdot)$ denotes *probability*.

5.1. Method - the OTM

¹ Let $\vec{\eta}$, do a certain number $e \cdot P$ of instructions $\eta_{[\cdot]}$ (of the program $\vec{\eta}$) beginning from the initial configuration $(\varepsilon, s_0, \vec{\xi})$. Let, e.g., $P = 2l + 1, l = \lfloor \vec{\xi}, e \geq 1; e \in \mathbb{N}$ be the number of the *stage* (for each stage we write, in a **programmer style**, $e := e + 1$)

- The step ¹ generates the table of nine-partite structures. Its length of $e \cdot P$ rows,¹³

$$\left[p P q P s_i x_k s_j y_l DP[\vec{\sigma}_{s_{[\cdot]}} \vec{\rho}] PC PP(\varepsilon[\sigma, s_{[\cdot]}, \rho] \varepsilon) P g P \vec{\sigma}_{s_{[\cdot]}} \vec{\rho} P m \right]_{p=1}^{p=e \cdot P}
 \tag{61}$$

where the denotation used is

- C is the number of the *configuration* $(\vec{\sigma}_{s_{[\cdot]}} \vec{\rho})$
- g is the number of the *configuration type* $(\varepsilon(\sigma, s_{[\cdot]}, \rho) \varepsilon)$
- m is the number of the *general configuration type* $G, (\vec{\sigma}_{s_{[\cdot]}} \vec{\rho})$

² In the table (¹) we are seeking two successive blocks of rows limited by those rows having the identical values in the columns (the identical six-partite structures)

$$\left[q P s_i x_k s_j y_l DP(\varepsilon[\sigma, s_{[\cdot]}, \rho] \varepsilon) P g P \vec{\sigma}_{s_{[\cdot]}} \vec{\rho} P m \right]
 \tag{62}$$

Thus, we are seeking for (the sequence of) the *three rows* being identical in those columns while the last row of the first block is the first row of the second block [this second ends by the third row (identical in the six columns considered)]. The numbers of these *separating* rows, the first,

¹³ The symbols '[' and ']' in the tables denote the range of the input (its limits) and, also, the 'operating space' for the CU_{TM} in each step.

the second and the third row are the numbers of steps $p_{[.]}, p_{[..]}$ and $p_{[...]}$ of the observed computing process. These rows are separated by numbers

$$\Delta p_{[.]} \triangleq p_{[.]} - p_{[.]} - 1 \geq 0, \quad \Delta p_{[..]} \triangleq p_{[..]} - p_{[.]} - 1 \geq 0 \tag{63}$$

of rows lying between them. (They can follow immediately, $\Delta p_{[.]}=0, \Delta p_{[..]}=0$.)

³ If the three separating rows are **not found** within the given stage e ($\overset{2}{-}$), we start the computing process driven by the program $\vec{\eta}$, and its tracing, from the beginning $[(\varepsilon, s_0, \vec{\xi}), (\overset{1}{-})]$, and let it run $e \cdot P$ steps, where $e := e + 1$ is set down.

⁴ If those three separating rows are **found** within the given stage e ($\overset{2}{-}$) (those two blocks covering the rows $p_{[.]}, p_{[...]}$ where $p_{[...]} = p_{[..]} + \Delta p_{[..]} + 1$ and $p_{[..]} = p_{[.]} + \Delta p_{[.]} + 1$) we are checking both the two blocks, each of them from its beginning ($p_{[.]}$ or $p_{[...]}$ respectively) till its end ($p_{[..]}$ or $p_{[...]}$ respectively), seeking the rows with the identical values within their six columns (62),¹⁴

$$\left[q P s_i x_k s_j y_l DP(\varepsilon [\sigma, s_{[.]}, \rho] \varepsilon) P g P \vec{\sigma}_{s_{[.]}} \vec{\rho} P m \right]$$

⁵ Two or more such **identical six-partite structures** on the successive row positions (denoted by symbols z and m) are **considered as the only one row**

$$\begin{aligned}
 & 1) \quad z_{[.]}, \dots, z_{[.]}, p_{[.]} = z_{[.]} \leq z_{[.]}, m = \overset{\text{Def}}{1} = m_{[.]} \\
 & 2) \quad z_{[.]}, \dots, z_{[.]}, z_{[.]} > z_{[.]}, z_{[.]} \geq z_{[.]}, m = 2 \\
 & 3) \quad z_{[.]}, \dots, z_{[.]}, z_{[.]} > z_{[.]}, z_{[.]} \geq z_{[.]}, m = 3 \\
 & \dots\dots\dots \\
 & \dots\dots) \quad z_{[.]}, \dots, z_{[.]}, p_{[.]} = z_{[.]} \geq z_{[.]}, m = \dots\dots \overset{\text{Def}}{=} m_{[.]}
 \end{aligned} \tag{64}$$

(the first of them is considered only).¹⁵

⁶ We check whether, by this way, two new **identical successive blocks** of unique six-partite structures are created [the first block is between the (newly numbered) rows $m_{[.]} \div m_{[..]}$ and the second is (newly) between the rows $m_{[..]} \div m_{[...]}$].

Their lengths are $\Delta m_{[.]}$ and $\Delta m_{[..]}$ [for $\Delta m_{[..]}: m_{[.]} := (m_{[..]})_{\text{mod}(m_{[..]}-1)}$]

¹⁴ If these blocks are identical the infinite cycle is discovered, but we continue uniformly with ⁵.

¹⁵ The situation in the first block is described only by (64).

$$\Delta m_{[.]} = m_{[.]} - m_{[.]} + 1 \leq \Delta p_{[.]}, \quad \Delta m_{[..]} = m_{[..]} - m_{[.]} + 1 \leq \Delta p_{[..]}. \quad (65)$$

If **NOT** - we continue with ², $p^* = p_{[...]}$; ¹ when $e \cdot P$ is exhausted, $e := e + 1$;

If **YES** - the distances ($\Delta m_{[.]}[-2]$ and $\Delta m_{[..]}[-2]$) between the marginal rows of the new blocks (⁵) are constant, $\Delta m_{[.]} = \Delta m_{[..]}$,

[the distances are counted in the number m of the unique six-partite structures (⁵), the last one of the first block is the first one of the second block], $\Delta m_{[.]} = \Delta m_{[..]}$.

⁷ If the distance $\Delta p_{[.]}$ of the first and the last row of the first block (²) is less or equal to the distance $\Delta p_{[..]}$ of the first and last row of the second blok $\Delta p_{[.]} \leq \Delta p_{[..]}$, we have discovered the infinite cycle driven by the program $\vec{\eta}$ [now the distances are counted in the number of steps p].

We continue further but **within the first block and with the $\Delta m_{[.]}$ only.**

From each unique six-partite structures (⁵, ⁶) the three columns

$$[qPGPm] \quad (66)$$

are now taken only, being interpreted as the rules of a regular grammar (accepted by a finite automaton)

$$S_{q^*} \rightarrow G_m S_{q_m} \quad (67)$$

$$q^* \in \{0\} \dot{\cup} \{1, \dots, |\vec{\eta}|\}, \quad q_m \in \{1, \dots, |\vec{\eta}|\}, \quad m = m_{[.]}, \dots, m_{[.]} - 1$$

with the set S of nonterminal symbols and the set T' of terminal symbols,

$$S_* \stackrel{\text{Def}}{=} \{S_0\} \dot{\cup} \{S_{q_m}\}_{m=m_{[.]}-1}^{m_{[.]}, \quad \text{card } S_* = n, \quad T' \stackrel{\text{Def}}{=} \{G_{q_m}\}_{m=m_{[.]}}^{m_{[.]}-1} \quad (68)$$

having the starting nonterminal symbol $S_0 \in S$.¹⁶

¹⁶ ^{7a} If the distance $\Delta p_{[.]}$ of the first and the last row of the first block (⁵ is greater than the distance $\Delta p_{[..]}$ of the first and last row of the second blok, $\Delta p_{[.]} > \Delta p_{[..]}$ we have discovered the finite cycle driven by the program $\vec{\eta}$.

⁸ For the first block of the unique six-partite structures (⁵, ⁶, ⁷) the sequence of rules of a regular grammar is being generated (accepted by the finite automaton with the states $S_0, S_{q_{m_{[.]}}}, \dots$)

$$\begin{aligned}
 S_0 &\rightarrow G_{m_{[.]}} S_{q_{m_{[.]}}} \\
 S_{q_{m_{[.]}}} &\rightarrow G_{m_{[.] + 1} } S_{q_{m_{[.] + 1} }} \\
 \dots &\dots \dots \\
 S_{q_{m_{[.] - 2} }} &\rightarrow G_{m_{[.] - 1} } S_{q_{m_{[.] - 1} }} \\
 S_{q_{m_{[.] - 2} }} &\rightarrow \bar{\varepsilon} S_0, \left[\text{or } S_{q_{m_{[.] - 1} }} \rightarrow \varepsilon S_{HALT}, S_{HALT} \notin \mathbf{S} \right]
 \end{aligned} \tag{69}$$

where the denotation $S_{q_{m_{[.]}} \triangleq S_{z_{[.]}, \dots, z_{[.]}} S_{q_{m_{[.] + 1}} \triangleq S_{z_{[.]}, \dots, z_{[.]}} \dots S_{q_{m_{[.]}} \triangleq S_{z_{[.]}, \dots, z_{[.]}}$ is used.

We have described the activity of a *finite automaton* which accepts the infinite regular language of the general configurations types (of the configurations of the observed machine *TM*). They are the words of the infinite length and having the form

$$\left[(\vec{\sigma}_m, s_m, \vec{\rho}_m)_{m=m_{[.]}}^{m_{[.] - 1} } \right]^+ = [\bar{\mathbf{X}}]^+ \left[\triangleq [\mathbf{d}(TM^*)]^+ \right] \tag{70}$$

Yet after the *second round* of the observed *TM* through the infinite cycle has been finished the **Pumping Lemma** is usable and valid for the length L of the relevant word of this infinite language, $[\text{cardS} \leq L < 2 \cdot \text{cardS}]$.

We can generate the status (the signal) S_{HALT} to halt the whole machine *OTM* and the *TM* consequently.

We can say, briefly, that: If such the three identical bi-partite structures $[\eta, (\vec{\sigma}, s, \vec{\rho})]$, following each other, exist that for their distances $\Delta p_{[.]}$ (measured by the number of steps of the observed process) between the first and the second bi-partite structure and $\Delta p_{[.]}$ between the second and the third bi-partite structure it is valid that $\Delta p_{[.]} \leq \Delta p_{[.]}$ the observed *TM* is going through the infinite cycle.

The expression (70) means that we have discovered, within the dynamical system $\mathcal{O}\mathcal{O}''$, the dynamical subsystem $\mathcal{O}^* \mathcal{O}' (\equiv TM^*)$ which is in a *limit cycle*. It means the thermodynamic equilibrium within the double cycle $\mathcal{O}^* \mathcal{O}'$, thus for its temperatures it is valid that $T_W^* = T_W^*$,

$T_0^* = T_0^*$; for sequences of the general configuration types $(G_{[\cdot]})$, \vec{X}_p and \vec{X}_{p+1} from (59), (60), it is valid, for certain $p \geq p^* \geq p^0 \geq 1$, that

$$H(\vec{X}_p) - H(\vec{X}_p | \vec{X}_{p+1}) = 0 \text{ where} \tag{71}$$

$$\vec{X}_p = (X_{p+1}, \dots, X_{2p+2-p^*}) = (X_{p^*}, \dots, X_p)$$

which represents the *zero change of the information and thermodynamic entropy* within the working medium of a reversible Carnot Cycle; for the sequences \vec{X}_p and \vec{X}_{p+1} of the observed TM 's configurations $X_{[\cdot]}$ from (47) ($TM \equiv O$) it is valid that

$$H(\vec{X}_p) - H(\vec{X}_p | \vec{X}_{p+1}) = H(\vec{Y}_p) > 0, \vec{X}_{[\cdot]} \triangleq X_{p^*}, \dots, X_{[\cdot]}, p^* \geq 1 \tag{72}$$

which represents the *nonzero output and, also, the growth of the thermodynamic and information entropy* within the whole isolated system in which that reversible Carnot Cycle is running, see [6, 8]. Generally, any Carnot Cycle is, under its construction draft, the infinite cycle, and, thus, both the relations (71) and (72) represent the *information thermodynamic criterion for the infinite cycle existence*.

The following section gives the examples of this method.

6. Examples

Example I

$$\vec{\eta} = (\eta_1, \eta_2, \eta_3, \eta_4);$$

$$\eta_1 = (s_0, I, s_0, I, R)$$

$$\eta_2 = (s_0, B, s_1, I, L) \tag{73}$$

$$\eta_3 = (s_1, I, s_1, I, L)$$

$$\eta_4 = (s_1, B, s_0, B, R)$$

This program *conserves* the given string $\vec{\xi}$

$$\vec{\xi} = [IIII...I] \text{ resp. } B[IIII...I]B \text{ resp. } B[IIII...I]B \tag{74}$$

or **BIB** respectively. Let the input $\vec{\xi} = IIII$ be given. From the table Tab. 1 it is obvious that

$$\begin{aligned}
 p_{[.]} &= 1, p_{[..]} = 13, p_{[...]} = 25, m_{[.]} = 1, m_{[..]} = 7 \text{ and thus} \\
 \Delta p_{[.]} &= 11, \Delta p_{[..]} = 13, \Delta p_{[...]} = 15, \Delta m_{[.]} = \Delta m_{[..]} = \Delta m_{[...]} = \dots = 7
 \end{aligned}
 \tag{75}$$

Then we can write down the regular grammar of the (regular) language of the general configuration given by the computing process driven by the program $\vec{\eta}$,

$$\begin{aligned}
 S_0 &\rightarrow \overline{B} s_0 \overline{IB} S_1 \\
 S_1 &\rightarrow \overline{BI} s_0 \overline{IB} S_{2345} \\
 S_{2345} &\rightarrow \overline{BI} s_0 \overline{B} S_6 \\
 S_6 &\rightarrow \overline{BI} s_1 \overline{IB} S_{78910} \\
 S_{78910} &\rightarrow \overline{B} s_1 \overline{IB} S_{11} \\
 S_{11} &\rightarrow \overline{B} s_1 \overline{BIB} S_{12} \\
 S_{12} &\rightarrow \vec{\varepsilon} S_0
 \end{aligned}
 \tag{76}$$

This grammar is with the set S_* of its nonterminal symbols,

$$S_* = \{S_0, S_1, S_{2345}, S_6, S_{78910}, S_{11}, S_{12}\}, \text{ card } S_* = 7
 \tag{77}$$

(let us remember (66)-(69) and that $S_1 \triangleq S_{q_1'}$, $S_2 \triangleq S_{q_2'}$, $S_6 \triangleq S_{q_3'}$, $S_7 \triangleq S_{q_4'}$, $S_{11} \triangleq S_{q_5'}$, $S_{12} = S_{q_6'}$)

and generates (the computing process described by it generates) the infinite word $[\overline{X}]^+$,

$$[\overline{X}]^+ = [\overline{B} s_0 \overline{IB}, \overline{BI} s_0 \overline{IB}, \overline{BI} s_0 \overline{B}, \overline{BI} s_1 \overline{IB}, \overline{B} s_1 \overline{IB}, \overline{B} s_1 \overline{BIB}]^+ = [d(TM^*)]^+
 \tag{78}$$

For the generation of the signal which halts the whole combined machine we can add and use the rule

$$S_{12} \rightarrow \varepsilon S_{HALT}, S_{HALT} \notin S_*
 \tag{79}$$

After the second round through the indicated infinite cycle the word of the general configuration types of the length $l=12$ is generated out and, thus, the Pumping Lemma it is valid,

$$\begin{aligned}
 \text{card } S_* \leq l < 2 \cdot (\text{card } S_*), \quad \text{card } S_* = 7 \\
 7 \leq 12 < 2 \cdot 7
 \end{aligned}
 \tag{80}$$

See the following table.

p ####	q ####	Instruction	Configuration	C ####	Config. Type	g ####	General Config. Type G	m ####
1	1	$s_0 I s_0 I R$	$\varepsilon B[s_0 I I I I] B \varepsilon$	1	$\varepsilon B[s_0 I] B \varepsilon$	1	$\vec{B} s_0 \vec{I B}$	1
2	1	$s_0 I s_0 I R$	$\varepsilon B[I s_0 I I I] B \varepsilon$	2	$\varepsilon B[I s_0 I] B \varepsilon$	2	$\vec{B I} s_0 \vec{I B}$	2
3	1	$s_0 I s_0 I R$	$\varepsilon B[I I s_0 I I] B \varepsilon$	3	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
4	1	$s_0 I s_0 I R$	$\varepsilon B[I I I s_0 I I] B \varepsilon$	4	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
5	1	$s_0 I s_0 I R$	$\varepsilon B[I I I I s_0 I] B \varepsilon$	5	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
6	2	$s_0 B s_1 B L$	$\varepsilon B[I I I I s_0] B \varepsilon$	6	$\varepsilon B[I s_0] B \varepsilon$	3	$\vec{B I} s_0 \vec{B}$	3
7	3	$s_1 I s_1 I L$	$\varepsilon B[I I I I s_1 I B] \varepsilon$	7	$\varepsilon B[I s_1 I B] \varepsilon$	4	$\vec{B I} s_1 \vec{I B}$	4
8	3	$s_1 I s_1 I L$	$\varepsilon B[I I I s_1 I I B] \varepsilon$	8	$\varepsilon B[I s_1 I B] \varepsilon$	4	4
9	3	$s_1 I s_1 I L$	$\varepsilon B[I I s_1 I I I B] \varepsilon$	9	$\varepsilon B[I s_1 I B] \varepsilon$	4	4
10	3	$s_1 I s_1 I L$	$\varepsilon B[I s_1 I I I I B] \varepsilon$	10	$\varepsilon B[I s_1 I B] \varepsilon$	4	4
11	3	$s_1 I s_1 I L$	$\varepsilon B[s_1 I I I I I B] \varepsilon$	11	$\varepsilon B[s_1 I B] \varepsilon$	5	$\vec{B} s_1 \vec{I B}$	5
12	4	$s_1 B s_0 B R$	$\varepsilon[s_1 B I I I I I B] \varepsilon$	12	$\varepsilon[s_1 B I B] \varepsilon$	6	$\vec{B} s_1 \vec{B I B}$	6
13	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I B] \varepsilon$	1'	$\varepsilon[B s_0 I B] \varepsilon$	1	$\vec{B} s_0 \vec{I B}$	7, 1
14	1	$s_0 I s_0 I R$	$\varepsilon[B I s_0 I I I I B] \varepsilon$	2'	$\varepsilon[B I s_0 I B] \varepsilon$	2	$\vec{B I} s_0 \vec{I B}$	2
15	1	$s_0 I s_0 I R$	$\varepsilon[B I I s_0 I I I B] \varepsilon$	3'	$\varepsilon[B I s_0 I B] \varepsilon$	2	2
16	1	$s_0 I s_0 I R$	$\varepsilon[B I I I s_0 I I B] \varepsilon$	4'	$\varepsilon[B I s_0 I B] \varepsilon$	2	2
17	1	$s_0 I s_0 I R$	$\varepsilon[B I I I I s_0 I B] \varepsilon$	5'	$\varepsilon[B I s_0 I B] \varepsilon$	2	2
18	2	$s_0 B s_1 B L$	$\varepsilon[B I I I I s_0 B] \varepsilon$	6'	$\varepsilon[B I s_0 B] \varepsilon$	3	$\vec{B I} s_0 \vec{B}$	3
19	3	$s_1 I s_1 I L$	$\varepsilon[B I I I I s_1 I B] \varepsilon$	7'	$\varepsilon[B I s_1 I B] \varepsilon$	4	$\vec{B I} s_1 \vec{I B}$	4
20	3	$s_1 I s_1 I L$	$\varepsilon[B I I I s_1 I I B] \varepsilon$	8'	$\varepsilon[B I s_1 I B] \varepsilon$	4	4
21	3	$s_1 I s_1 I L$	$\varepsilon[B I I s_1 I I I B] \varepsilon$	9'	$\varepsilon[B I s_1 I B] \varepsilon$	4	4
22	3	$s_1 I s_1 I L$	$\varepsilon[B I s_1 I I I I B] \varepsilon$	10'	$\varepsilon[B I s_1 I B] \varepsilon$	4	4
23	3	$s_1 I s_1 I L$	$\varepsilon[B s_1 I I I I I B] \varepsilon$	11'	$\varepsilon[B s_1 I B] \varepsilon$	5	$\vec{B} s_1 \vec{I B}$	5
24	4	$s_1 B s_0 B R$	$\varepsilon[s_1 B I I I I I B] \varepsilon$	12'	$\varepsilon[s_1 B I B] \varepsilon$	6	$\vec{B} s_1 \vec{B I B}$	6
25	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I B] \varepsilon$	1''	$\varepsilon[B s_0 I B] \varepsilon$	1	$\vec{B} s_0 \vec{I B}$	7, 1
26	2''	2	2
..
36	12''	6	6
37	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I B] \varepsilon$	1'''	$\varepsilon[B s_0 I B] \varepsilon$	1	$\vec{B} s_0 \vec{I B}$	1
38	2'''	2	2
..
48	12'''	6	6
49	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I B] \varepsilon$	1''''	$\varepsilon[B s_0 I B] \varepsilon$	1	$\vec{B} s_0 \vec{I B}$	1

Table 1. Tracing and staging for Example I

[We 'idle' for the same number of clicks Δp in each of the blocks Δm . The sequence of Δp is (11, 11, 11, ...), the sequence of Δm is (5, 5, 5, ...). The sequence for p is (13, 13, 13, ...) and for m is (7, 7, 7, ...).]

After the *observed (sub)machine* has entered into the infinite cycle, which takes place in a finite time, and has gone through this cycle twice it is halted by the signal from the observing machine.

Example II

$$\begin{aligned}
 \vec{\eta} &= (\eta_1, \eta_2, \eta_3, \eta_4) \\
 \eta_1 &= (s_0, I, s_0, I, R) \\
 \eta_2 &= (s_0, B, s_1, I, L) \\
 \eta_3 &= (s_1, I, s_1, I, L) \\
 \eta_4 &= (s_1, B, s_0, B, R)
 \end{aligned}
 \tag{81}$$

This program generates the expanding sequence

$$[IIII] \text{ resp. } B[IIII...I]B, \dots, B[IIII...I...]B, \dots
 \tag{82}$$

or **BIB** respectively. Let the input $\vec{\xi} = IIII$ be given. From the following table Tab. 2 it is obvious that

$$\begin{aligned}
 p_{[.]} &= 1, p_{[..]} = 13, p_{[...]} = 27, m_{[.]} = 1, m_{[..]} = 7 \text{ and thus} \\
 \Delta p_{[.]} &= 11, \Delta p_{[..]} = 13, \Delta p_{[...]} = 15, \Delta m_{[.]} = \Delta m_{[..]} = \Delta m_{[...]} = \dots = 7
 \end{aligned}
 \tag{83}$$

and we can write down the regular grammar (of the regular) language of the general configuration types generated by the computing process driven by $\vec{\eta}$

$$\begin{aligned}
 S_0 &\rightarrow \overline{B} s_0 \overline{IB} S_1 \\
 S_1 &\rightarrow \overline{BI} s_0 \overline{IB} S_{2345} \\
 S_{2345} &\rightarrow \overline{BI} s_0 \overline{B} S_6 \\
 S_6 &\rightarrow \overline{BI} s_1 \overline{IB} S_{78910} \\
 S_{78910} &\rightarrow \overline{B} s_1 \overline{IB} S_{11} \\
 S_{11} &\rightarrow \overline{B} s_1 \overline{BIB} S_{12} \\
 S_{12} &\rightarrow \vec{\epsilon} S_0
 \end{aligned}
 \tag{84}$$

where, after the relations (66)-(69), $S_1 \triangleq S_{q_1}$, $S_2 \triangleq S_{q_2}$, $S_6 \triangleq S_{q_3}$, $S_7 \triangleq S_{q_4}$, $S_{11} \triangleq S_{q_5}$, $S_{12} = S_{q_6}$.

This grammar is with the set of nonterminal symbols

$$S_* = \{S_0, S_1, S_{2345}, S_6, S_{78910}, S_{11}, S_{12}\}, \text{card } S_* = 7 \quad (85)$$

and generates (the computing process described by it generates) the infinite word $[\vec{X}]^+$,

$$[\vec{X}]^+ = [\overline{B} s_0, \overline{IB}, \overline{BI} s_0, \overline{IB}, \overline{BI} s_0, \overline{B}, \overline{BI} s_1, \overline{IB}, \overline{B} s_1, \overline{IB}, \overline{B} s_1, \overline{BIB}]^+ = [\mathbf{d}(TM^*)]^+ \quad (86)$$

For the generation of the signal which halts the whole combined machine we can add and use the rule

$$S_{12} \rightarrow \varepsilon S_{\text{HALT}}, S_{\text{HALT}} \notin S_* \quad (87)$$

After the second round through the indicated infinite cycle the word of the general configuration types of the length $l=12$ is generated and, thus, the Pumping Lemma it is valid,

$$\begin{aligned} \text{card } S_* \leq l < 2 \cdot (\text{card } S_*), \quad \text{card } S_* = 7 \\ 7 \leq 12 < 2 \cdot 7 \end{aligned} \quad (88)$$

See the following table.

We are going through the seven states repeatedly and each such a pass lasts longer than the previous one.

We 'idle' in such each pass a longer time (for the growing number of clicks p of the CU_{TM}).

[The growing sequence of Δp is (11, 13, 15,...) is evidenced while the sequence of Δm is (5,5,5,...).

The sequence for p is (13, 15, 17,...) and for m is (7, 7, 7,...).]

After the *observed (sub)machine* has entered into the infinite cycle, which occurs in a finite time, and has gone through this cycle twice *it is halted by the signal from the observing machine*.

p ####	q ####	Instruction	Configuration	C ####	Config. Type	g ####	General Config. Type G	m ####
1	1	$s_0 I s_0 I R$	$\varepsilon B[s_0 I I I I] B \varepsilon$	1	$\varepsilon B[s_0 I] B \varepsilon$	1	$\vec{B}_{s_0} \vec{I B}$	1
2	1	$s_0 I s_0 I R$	$\varepsilon B[I s_0 I I I] B \varepsilon$	2	$\varepsilon B[I s_0 I] B \varepsilon$	2	$\vec{B I}_{s_0} \vec{I B}$	2
3	1	$s_0 I s_0 I R$	$\varepsilon B[I I s_0 I I] B \varepsilon$	3	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
4	1	$s_0 I s_0 I R$	$\varepsilon B[I I I s_0 I I] B \varepsilon$	4	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
5	1	$s_0 I s_0 I R$	$\varepsilon B[I I I I s_0 I] B \varepsilon$	5	$\varepsilon B[I s_0 I] B \varepsilon$	2	2
6	2	$s_0 B s_1 I L$	$\varepsilon B[I I I I s_0 B] \varepsilon$	6	$\varepsilon B[I s_0 B] \varepsilon$	3	$\vec{B I}_{s_0} \vec{B}$	3
7	3	$s_1 I s_1 I L$	$\varepsilon B[I I I I s_1 I I] \varepsilon$	7	$\varepsilon B[I s_1 I] \varepsilon$	4	$\vec{B I}_{s_1} \vec{I B}$	4
8	3	$s_1 I s_1 I L$	$\varepsilon B[I I I s_1 I I I] \varepsilon$	8	$\varepsilon B[I s_1 I] \varepsilon$	4	4
9	3	$s_1 I s_1 I L$	$\varepsilon B[I I s_1 I I I I] \varepsilon$	9	$\varepsilon B[I s_1 I] \varepsilon$	4	4
10	3	$s_1 I s_1 I L$	$\varepsilon B[I s_1 I I I I I] \varepsilon$	10	$\varepsilon B[I s_1 I] \varepsilon$	4	4
11	3	$s_1 I s_1 I L$	$\varepsilon B[s_1 I I I I I I] \varepsilon$	11	$\varepsilon B[s_1 I] \varepsilon$	5	$\vec{B}_{s_1} \vec{I B}$	5
12	4	$s_1 B s_0 B R$	$\varepsilon [s_1 B I I I I I I] \varepsilon$	12	$\varepsilon [s_1 B I] \varepsilon$	6	$\vec{B}_{s_1} \vec{B I B}$	6
13	1	$s_0 I s_0 I R$	$\varepsilon [B s_0 I I I I I I] B \varepsilon$	1'	$\varepsilon [B s_0 I] B \varepsilon$	1	$\vec{B}_{s_0} \vec{I B}$	7, 1
14	1	$s_0 I s_0 I R$	$\varepsilon [B I s_0 I I I I] B \varepsilon$	2'	$\varepsilon [B I s_0 I] B \varepsilon$	2	$\vec{B I}_{s_0} \vec{I B}$	2
15	1	$s_0 I s_0 I R$	$\varepsilon [B I I s_0 I I I] B \varepsilon$	3'	$\varepsilon [B I s_0 I] B \varepsilon$	2	2
16	1	$s_0 I s_0 I R$	$\varepsilon [B I I I s_0 I I] B \varepsilon$	4'	$\varepsilon [B I s_0 I] B \varepsilon$	2	2
17	1	$s_0 I s_0 I R$	$\varepsilon [B I I I I s_0 I I] B \varepsilon$	5'	$\varepsilon [B I s_0 I] B \varepsilon$	2	2
18	1	$s_0 I s_0 I R$	$\varepsilon [B I I I I I s_0 I] B \varepsilon$	5'	$\varepsilon [B I s_0 I] B \varepsilon$	2	2
19	2	$s_0 B s_1 I L$	$\varepsilon [B I I I I I s_0 B] \varepsilon$	6'	$\varepsilon [B I s_0 B] \varepsilon$	3	$\vec{B I}_{s_0} \vec{B}$	3
20	3	$s_1 I s_1 I L$	$\varepsilon [B I I I I I s_1 I I] \varepsilon$	7'	$\varepsilon [B I s_1 I] \varepsilon$	4	$\vec{B I}_{s_1} \vec{I B}$	4
21	3	$s_1 I s_1 I L$	$\varepsilon [B I I I I s_1 I I I] \varepsilon$	8'	$\varepsilon [B I s_1 I] \varepsilon$	4	4
22	3	$s_1 I s_1 I L$	$\varepsilon [B I I I s_1 I I I I] \varepsilon$	9'	$\varepsilon [B I s_1 I] \varepsilon$	4	4
23	3	$s_1 I s_1 I L$	$\varepsilon [B I I s_1 I I I I I] \varepsilon$	10'	$\varepsilon [B I s_1 I] \varepsilon$	4	4
24	3	$s_1 I s_1 I L$	$\varepsilon [B I s_1 I I I I I I] \varepsilon$	10'	$\varepsilon [B I s_1 I] \varepsilon$	4	4
25	3	$s_1 I s_1 I L$	$\varepsilon [B s_1 I I I I I I I] \varepsilon$	11'	$\varepsilon [B s_1 I] \varepsilon$	5	$\vec{B}_{s_1} \vec{I B}$	5
26	4	$s_1 B s_0 B R$	$\varepsilon [s_1 B I I I I I I I] \varepsilon$	12'	$\varepsilon [s_1 B I] \varepsilon$	6	$\vec{B}_{s_1} \vec{B I B}$	6
27	1	$s_0 I s_0 I R$	$\varepsilon [B s_0 I I I I I I I] B \varepsilon$	1''	$\varepsilon [B s_0 I] B \varepsilon$	1	$\vec{B}_{s_0} \vec{I B}$	7, 1

p ####	q ####	Instruction	Configuration	C ####	Config. Type	g ####	General Config. Type G	m ####
27	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I I]B\varepsilon$	1''	$\varepsilon[B s_0 I]B\varepsilon$	1	$\vec{B} s_0 \vec{I}B$	1
28	1	$s_0 I s_0 I R$	$\varepsilon[B I s_0 I I I I I]B\varepsilon$	2''	$\varepsilon[B I s_0 I]B\varepsilon$	2	$\vec{B} I s_0 \vec{I}B$	2
29	1	$s_0 I s_0 I R$	$\varepsilon[B I I s_0 I I I I]B\varepsilon$	3''	$\varepsilon[B I s_0 I]B\varepsilon$	2	2
30	1	$s_0 I s_0 I R$	$\varepsilon[B I I I s_0 I I I]B\varepsilon$	4''	$\varepsilon[B I s_0 I]B\varepsilon$	2	2
31	1	$s_0 I s_0 I R$	$\varepsilon[B I I I I s_0 I I I]B\varepsilon$	5''	$\varepsilon[B I s_0 I]B\varepsilon$	2	2
32	1	$s_0 I s_0 I R$	$\varepsilon[B I I I I I s_0 I I]B\varepsilon$	5''	$\varepsilon[B I s_0 I]B\varepsilon$	2	2
33	1	$s_0 I s_0 I R$	$\varepsilon[B I I I I I I s_0 I]B\varepsilon$	5''	$\varepsilon[B I s_0 I]B\varepsilon$	2	2
34	2	$s_0 B s_1 I L$	$\varepsilon[B I I I I I I s_0 B]\varepsilon$	6''	$\varepsilon[B I s_0 B]\varepsilon$	3	$\vec{B} I s_0 \vec{B}$	3
35	3	$s_1 I s_1 I L$	$\varepsilon[B I I I I I I s_1 I I]\varepsilon$	7''	$\varepsilon[B I s_1 I]\varepsilon$	4	$\vec{B} I s_1 \vec{I}B$	4
36	3	$s_1 I s_1 I L$	$\varepsilon[B I I I I I I s_1 I I I]\varepsilon$	8''	$\varepsilon[B I s_1 I]\varepsilon$	4	4
37	3	$s_1 I s_1 I L$	$\varepsilon[B I I I I I s_1 I I I I]\varepsilon$	9''	$\varepsilon[B I s_1 I]\varepsilon$	4	4
38	3	$s_1 I s_1 I L$	$\varepsilon[B I I I s_1 I I I I I]\varepsilon$	10''	$\varepsilon[B I s_1 I]\varepsilon$	4	4
39	3	$s_1 I s_1 I L$	$\varepsilon[B I I s_1 I I I I I I]\varepsilon$	10''	$\varepsilon[B I s_1 I]\varepsilon$	4	4
40	3	$s_1 I s_1 I L$	$\varepsilon[B I s_1 I I I I I I I]\varepsilon$	10''	$\varepsilon[B I s_1 I]\varepsilon$	4	4
41	3	$s_1 I s_1 I L$	$\varepsilon[B s_1 I I I I I I I I]\varepsilon$	11''	$\varepsilon[B s_1 I]\varepsilon$	5	$\vec{B} s_1 \vec{I}B$	5
42	4	$s_1 B s_0 B R$	$\varepsilon[s_1 B I I I I I I I]\varepsilon$	12''	$\varepsilon[s_1 B I]\varepsilon$	6	$\vec{B} s_1 \vec{B I}B$	6
43	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I I I I]B\varepsilon$	1'''	$\varepsilon[B s_0 I]B\varepsilon$	1	$\vec{B} s_0 \vec{I}B$	7, 1
44	2'''	2	2
..
60	12'''	6	6
61	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I I I I I]B\varepsilon$	1''''	$\varepsilon[B s_0 I]B\varepsilon$	1	$\vec{B} s_0 \vec{I}B$	7, 1
62	2''''	2	2
..
80	12''''	6	6
81	1	$s_0 I s_0 I R$	$\varepsilon[B s_0 I I I I I I I I I I]B\varepsilon$	1'''''	$\varepsilon[B s_0 I]B\varepsilon$	1	$\vec{B} s_0 \vec{I}B$	7, 1

Table 2. Tracing and staging for Example II

Example III

$$\begin{aligned}
 \vec{\eta} &= (\eta_1, \eta_2, \eta_3, \eta_4, \eta_5, \eta_6, \eta_7, \eta_8, \eta_9, \eta_{10}, \eta_{11}, \eta_{12}, \eta_{13}) \\
 \eta_1 &= (s_0, I, s_1, B, R) \\
 \eta_2 &= (s_1, I, s_1, I, R) \\
 \eta_3 &= (s_1, B, s_2, B, R) \\
 \eta_4 &= (s_2, I, s_2, I, R) \\
 \eta_5 &= (s_2, B, s_3, B, L) \\
 \eta_6 &= (s_3, I, s_4, B, L) \\
 \eta_7 &= (s_4, B, s_H, B, R) \\
 \eta_8 &= (s_4, I, s_5, I, L) \\
 \eta_9 &= (s_5, I, s_5, I, L) \\
 \eta_{10} &= (s_5, B, s_6, B, L) \\
 \eta_{11} &= (s_6, I, s_6, I, L) \\
 \eta_{12} &= (s_6, B, s_0, B, R) \\
 \eta_{13} &= (s_0, B, s_0, B, L)
 \end{aligned}
 \tag{89}$$

This program *generates*, or is figuring, the *difference* between the numbers 3 and 4 in this order,

$$3 - 4 \tag{90}$$

and *without checking the negativity* of the result. Thus it enters into the infinite cycle in which generates the left-direction expanding sequence

$$[\dots BBB \dots BIBBB] \text{ resp. } B [\dots BBB \dots BIBBB] B \cong \overline{BIB} \tag{91}$$

Let now the input word $\vec{\xi} = IIIIIII$ be given.

From the following table Tab 3. it is obvious that after the generation of the finite length's beginning part G, but of the infinite lengthy word of the general configurations numbered by values of m ,

$$\begin{aligned}
 \mathbf{G} = [& \begin{array}{cccccccccccccccc}
 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & 15 & 16 \\
 1 & 2 & 3 & 5 & 6 & 7 & 8 & 10 & 12 & 13 & 15 & 16 & & & & \\
 1 & 17 & 18 & 6 & 8 & 19 & 20 & 21 &] & & & & & & &
 \end{array}
 \end{aligned}
 \tag{92}$$

the part [22]⁺ follows. From the table it is obvious that the cycle is with the values

$$\begin{aligned} p_{[.]} &= 41, p_{[..]} = 42, p_{[...]} = 43 \\ m_{[.]} &= 41, m_{[..]} = 42 \end{aligned} \quad (93)$$

We can write down the regular grammar (of the regular) language of the general configuration types generated by the computing process driven by $\vec{\eta}$

$$\begin{aligned} S_0 &\rightarrow \mathbf{G} S_{1\dots 40} \\ S_{1\dots 40} &\rightarrow \overline{\mathbf{B}} s_0 \overline{\mathbf{BIB}} S_{414243\dots} \\ S_{414243\dots} &\rightarrow \overline{\mathbf{B}} s_0 \overline{\mathbf{BIB}} S_{414243\dots} \end{aligned} \quad (94)$$

This grammar is with the set of nonterminal symbols

$$\mathbf{S}_* = \{S_0, S_{1\dots 40}, S_{414243\dots}, S_6, S_{78910}, S_{11}, S_{12}\}, \text{card } \mathbf{S}_* = 3 \quad (95)$$

and generates the infinite word

$$\left[\mathbf{G}[22]^+ \right] \quad (96)$$

After the second round through the indicated infinite cycle the subword of the length $l = 12$ is generated out,

$$\left[\overline{\mathbf{B}} s_0 \overline{\mathbf{BIB}}, \overline{\mathbf{B}} s_0 \overline{\mathbf{BIB}}, \overline{\mathbf{B}} s_0 \overline{\mathbf{BIB}} \right] \quad (97)$$

and thus, the Pumping Lemma it is valid,

$$\begin{aligned} \text{card } \mathbf{S}_* \leq l < 2 \cdot (\text{card } \mathbf{S}_*), \quad \text{card } \mathbf{S}_* = 3 \\ 3 \leq 3 < 2 \cdot 3 \end{aligned} \quad (98)$$

For the generation of the signal which halts the whole combined machine we can use (or add) the rule

$$S_{414243\dots} \rightarrow \varepsilon S_{\text{HALT}}, S_{\text{HALT}} \notin \mathbf{S}_* \quad (99)$$

Here again it is valid that for the finite number of tracing steps and, for each is lasting the finitely long time, both halt states of our double-Turing Machine, which is the Turing Machine, too, occur *in a finite time*.

p ####	q ####	Instruction	Configuration	C ####	Config. Type	g ####	General Config. Type G	m ####
1	1	$s_0 I s_1 I R$	$\varepsilon B[s_0 IIIIBIII]B\varepsilon$	1	$\varepsilon B[s_0 IBI]B\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1
2	2	$s_1 I s_1 I R$	$\varepsilon B[s_1 IIBIII]B\varepsilon$	2	$\varepsilon B[s_1 IBI]B\varepsilon$	2	$\vec{B} s_1 \vec{IB}$	2
3	2	$s_1 I s_1 I R$	$\varepsilon B[I s_1 IBIII]B\varepsilon$	3	$\varepsilon B[I s_1 IBI]B\varepsilon$	3	$\vec{BI} s_1 \vec{IB}$	3
4	3	$s_1 B s_2 B R$	$\varepsilon B[II s_1 BIII]B\varepsilon$	4	$\varepsilon B[I s_1 BI]B\varepsilon$	4	$\vec{BI} s_1 \vec{BIB}$	4
5	4	$s_2 I s_2 I R$	$\varepsilon B[IIB s_2 III]B\varepsilon$	5	$\varepsilon B[IB s_2 I]B\varepsilon$	5	$\vec{BIB} s_2 \vec{IB}$	5
6	4	$s_2 I s_2 I R$	$\varepsilon B[IIBI s_2 II]B\varepsilon$	6	$\varepsilon B[IBI s_2 I]B\varepsilon$	6	$\vec{BI} s_2 \vec{IB}$	6
7	4	$s_2 I s_2 I R$	$\varepsilon B[IIBII s_2 I]B\varepsilon$	7	$\varepsilon B[IBI s_2 I]B\varepsilon$	6	6
8	5	$s_2 B s_3 B L$	$\varepsilon B[IIBIII s_2]B\varepsilon$	8	$\varepsilon B[IBIBI s_2]B\varepsilon$	7	$\vec{BI} s_2 \vec{B}$	7
9	6	$s_3 I s_4 B L$	$\varepsilon B[IIBII s_3 IB]\varepsilon$	9	$\varepsilon B[IBI s_3 IB]\varepsilon$	8	$\vec{BI} s_3 \vec{IB}$	8
10	6	$s_3 I s_4 B L$	$\varepsilon B[IIBIII s_3 BB]\varepsilon$	10	$\varepsilon B[IBI s_3 B]\varepsilon$	9	$\vec{BI} s_3 \vec{B}$	9
11	8	$s_4 I s_5 I L$	$\varepsilon B[IIBII s_4 IBB]\varepsilon$	11	$\varepsilon B[IBI s_4 IB]\varepsilon$	10	$\vec{BI} s_4 \vec{IB}$	10
12	9	$s_5 I s_5 I L$	$\varepsilon B[IIBI s_5 IIBB]\varepsilon$	12	$\varepsilon B[IBI s_5 IB]\varepsilon$	11	$\vec{BI} s_5 \vec{IB}$	11
13	9	$s_5 I s_5 I L$	$\varepsilon B[IIB s_5 IIIBB]\varepsilon$	13	$\varepsilon B[IB s_5 IB]\varepsilon$	12	$\vec{BIB} s_5 \vec{IB}$	12
14	10	$s_5 B s_6 B L$	$\varepsilon B[II s_5 BIIIBB]\varepsilon$	14	$\varepsilon B[I s_5 BIB]\varepsilon$	13	$\vec{BI} s_5 \vec{BIB}$	13
15	11	$s_6 I s_6 I L$	$\varepsilon B[I s_6 IBIIIBB]\varepsilon$	15	$\varepsilon B[I s_6 IBIB]\varepsilon$	14	$\vec{BI} s_6 \vec{IB}$	14
16	11	$s_6 I s_6 I L$	$\varepsilon B[s_6 IIBIIIBB]\varepsilon$	16	$\varepsilon B[s_6 IBIB]\varepsilon$	15	$\vec{B} s_6 \vec{IB}$	15
17	12	$s_6 B s_0 B R$	$\varepsilon[s_6 BIIBIIIBB]\varepsilon$	17	$\varepsilon[s_6 BIBIB]\varepsilon$	16	$\vec{B} s_6 \vec{BIB}$	16
18	1	$s_0 I s_1 I R$	$\varepsilon[B s_0 IIBIIIBB]\varepsilon$	1	$\varepsilon[B s_0 IBIB]\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1
19	2	$s_1 I s_1 I R$	$\varepsilon[B s_1 IBIIIBB]\varepsilon$	2'	$\varepsilon[B s_1 IBIB]\varepsilon$	2	$\vec{B} s_1 \vec{IB}$	2
20	3	$s_1 B s_2 B R$	$\varepsilon[B I s_1 BIIIBB]\varepsilon$	4'	$\varepsilon[B I s_1 BIB]\varepsilon$	3	$\vec{BI} s_1 \vec{BIB}$	3
21	4	$s_2 I s_2 I R$	$\varepsilon[B IB s_2 IIIBB]\varepsilon$	5'	$\varepsilon[B IB s_2 IB]\varepsilon$	5	$\vec{BIB} s_2 \vec{IB}$	5
22	4	$s_2 I s_2 I R$	$\varepsilon[B IBI s_2 IIBB]\varepsilon$	6'	$\varepsilon[B IBI s_2 IB]\varepsilon$	5	$\vec{BI} s_2 \vec{IB}$	6
23	4	$s_2 I s_2 I R$	$\varepsilon[B IBII s_2 IBB]\varepsilon$	7'	$\varepsilon[B IBI s_2 IB]\varepsilon$	6	6
24	5	$s_2 B s_3 B L$	$\varepsilon[B IBIII s_2 BB]\varepsilon$	8'	$\varepsilon[B IBI s_2 B]\varepsilon$	7	$\vec{BI} s_2 \vec{B}$	7
25	6	$s_3 I s_4 B L$	$\varepsilon[B IBII s_3 IBB]\varepsilon$	9'	$\varepsilon[B IBI s_3 IB]\varepsilon$	8	$\vec{BI} s_3 \vec{IB}$	8
26	8	$s_4 I s_5 I L$	$\varepsilon[B IBI s_4 IBBB]\varepsilon$	11'	$\varepsilon[B IBI s_4 IB]\varepsilon$	10	$\vec{BI} s_4 \vec{IB}$	10
27	9	$s_5 I s_5 I L$	$\varepsilon[B IB s_5 IIBBB]\varepsilon$	13'	$\varepsilon[B IB s_5 IB]\varepsilon$	12	$\vec{BIB} s_5 \vec{IB}$	12
28	10	$s_5 B s_6 B L$	$\varepsilon[B I s_5 BIIIBB]\varepsilon$	15'	$\varepsilon[B I s_5 BIB]\varepsilon$	13	$\vec{BI} s_5 \vec{BIB}$	13
29	11	$s_6 I s_6 I L$	$\varepsilon[B s_6 IBIIIBB]\varepsilon$	16'	$\varepsilon[B s_6 IBIB]\varepsilon$	15	$\vec{B} s_6 \vec{IB}$	15
30	12	$s_6 B s_0 B R$	$\varepsilon[s_6 BIBIIBBB]\varepsilon$	17'	$\varepsilon[s_6 BIBIB]\varepsilon$	16	$\vec{B} s_6 \vec{BIB}$	16
31	1	$s_0 I s_1 I R$	$\varepsilon[B s_0 IBIIIBB]\varepsilon$	1	$\varepsilon[B s_0 IBIB]\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1

p ####	q ####	Instruction	Configuration	C ####	Config. Type	g ####	General Config. Type G	m ####
31	1	$s_0 I s_1 I R$	$\varepsilon[B s_0 IBII BBB]\varepsilon$	1	$\varepsilon[B s_0 IBIB]\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1
32	3	$s_1 B s_2 B R$	$\varepsilon[BB s_1 BIIBBB]\varepsilon$	17	$\varepsilon[B s_1 BIB]\varepsilon$	17	$\vec{B} s_1 \vec{BIB}$	17
33	4	$s_2 I s_2 I R$	$\varepsilon[BBB s_2 IIBBB]\varepsilon$	18	$\varepsilon[B s_2 IB]\varepsilon$	18	$\vec{B} s_2 \vec{IB}$	18
34	4	$s_2 I s_2 I R$	$\varepsilon[BBBI s_2 IBBB]\varepsilon$	19	$\varepsilon[BI s_2 IB]\varepsilon$	19	$\vec{BI} s_2 \vec{IB}$	6
35	5	$s_2 B s_3 B L$	$\varepsilon[BBBII s_2 BBB]\varepsilon$	20	$\varepsilon[BI s_2 IB]\varepsilon$	19	6
36	6	$s_3 I s_4 B L$	$\varepsilon[BBBI s_3 IBBB]\varepsilon$	21	$\varepsilon[BI s_3 IB]\varepsilon$	20	$\vec{BI} s_3 \vec{IB}$	8
37	8	$s_4 I s_5 I L$	$\varepsilon[BBB s_4 IBBBB]\varepsilon$	22	$\varepsilon[B s_4 IB]\varepsilon$	21	$\vec{B} s_4 \vec{IB}$	19
38	10	$s_5 B s_6 B L$	$\varepsilon[BB s_5 BIBBB]\varepsilon$	23	$\varepsilon[B s_5 BIB]\varepsilon$	22	$\vec{B} s_5 \vec{IB}$	20
39	12	$s_6 B s_0 B R$	$\varepsilon[B s_6 BBIBBB]\varepsilon$	24	$\varepsilon[B s_6 BIB]\varepsilon$	23	$\vec{B} s_6 \vec{BIB}$	21
40	12	$s_6 B s_0 B R$	$\varepsilon[BB s_6 BIBBB]\varepsilon$	25	$\varepsilon[B s_6 BIB]\varepsilon$	23	$\vec{B} s_6 \vec{BIB}$	21
41	13	$s_0 B s_0 B L$	$\varepsilon[B s_0 BBIBBBB]\varepsilon$	26	$\varepsilon[B s_0 BIB]\varepsilon$	24	$\vec{B} s_0 \vec{BIB}$	22
42	13	$s_0 B s_0 B L$	$\varepsilon B[s_0 BBBIBBB]\varepsilon$	26'	$\varepsilon B[s_0 BIB]\varepsilon$	24	22
43	13	$s_0 B s_0 B L$	$\varepsilon B[s_0 BBBBIBBB]\varepsilon$	26''	$\varepsilon B[s_0 BIB]\varepsilon$	24	22
44	13	$s_0 B s_0 B L$	$\varepsilon B[s_0 BBBBIBBB]\varepsilon$	26'''	$\varepsilon B[s_0 BIB]\varepsilon$	24	22
45	13	26''''	24	22
..
..	13	24	22
....	13	$s_0 B s_0 B L$	$\varepsilon B[s_0 B...B...BIBBB]\varepsilon$	26'....'	$\varepsilon B[s_0 BIB]\varepsilon$	24	$\vec{B} s_0 \vec{BIB}$	22
..	13	24	22
..

Table 3. Tracing and staging for Example III

[The distance Δp of p for the general configuration numbered 22 is constant ($\Delta p=0$).]

Example IV

$$\begin{aligned}
 \vec{\eta} &= (\eta_1, \eta_2, \eta_3, \eta_4, \eta_5, \eta_6, \eta_7, \eta_8, \eta_9, \eta_{10}, \eta_{11}, \eta_{12}, \eta_{13}) \\
 \eta_1 &= (s_0, I, s_1, B, R) \\
 \eta_2 &= (s_1, I, s_1, I, R) \\
 \eta_3 &= (s_1, B, s_2, B, R) \\
 \eta_4 &= (s_2, I, s_2, I, R) \\
 \eta_5 &= (s_2, B, s_3, B, L)
 \end{aligned}
 \tag{100}$$

$$\begin{aligned}
 \eta_6 &= (s_3, I, s_4, B, L) \\
 \eta_7 &= (s_4, B, s_H, B, R) \\
 \eta_8 &= (s_4, I, s_5, I, L) \\
 \eta_9 &= (s_5, I, s_5, I, L) \\
 \eta_{10} &= (s_5, B, s_6, B, L) \\
 \eta_{11} &= (s_6, I, s_6, I, L) \\
 \eta_{12} &= (s_6, B, s_0, B, R) \\
 \eta_{13} &= (s_0, B, s_0, B, L)
 \end{aligned}$$

This program is figuring the difference 4-3 and the input is $\vec{\xi} = IIIBIII$.

From the following table follows that during our staging the whole double-machine halts itself. In the numbers m of the general configuration types the following word \mathbf{G} is generated,

$$\begin{aligned}
 \mathbf{G} &= [\mathbf{G}_1 \mathbf{G}_2], \quad l = 2 \\
 \mathbf{G}_1 &= [1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 10\ 11\ 12] \\
 \mathbf{G}_2 &= [1\ 2\ 3\ 5\ 13\ 14\ 15\ 16]
 \end{aligned} \tag{101}$$

We can write the regular grammar of the (regular) language of the general configuration types being generated by the process driven by $\vec{\eta}$,

$$\begin{aligned}
 S_0 &\rightarrow \mathbf{G}_1 S_{1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 10\ 11\ 12\ 13} \\
 S_{1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 10\ 11\ 12\ 13} &\rightarrow \mathbf{G}_2 S_{14\ 15\ 16\ 17\ 18\ 19\ 20\ 21}
 \end{aligned} \tag{102}$$

This grammar has the set \mathbf{S}_* of the nonterminal symbols,

$$\mathbf{S}_* = \{S_0, S_{1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 10\ 11\ 12\ 13}, S_{14\ 15\ 16\ 17\ 18\ 19\ 20\ 21} \equiv S_{HALT}\}, \quad \text{card } \mathbf{S}_* = 3 \tag{103}$$

The whole *staging is ended by the natural end* of the process in the *observed machine*.

[All computing and tracing is stopped by the *end of the observed process*.]

All the previous examples have shown that after the finite number of steps, each is interpreted in a finite time and thus we need the finite time only, the signal halting the whole double machine is generated in the *all cases*. The reason for generating the *halting signal* and as such, the recognition the finite or infinite cycle in the observed machine is quite visible from all our tracing tables. Our double machine is the Turing Machine, too.

p	q	Instruction	Configuration	C	Config. Type	g	General Config. Type G	m
####	####			####		####		####
1	1	$s_0 I s_1 B R$	$\varepsilon[s_0 III BII]B\varepsilon$	1	$\varepsilon[s_0 IBI]\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1
2	2	$s_1 I s_1 I R$	$\varepsilon[B s_1 IIBII]B\varepsilon$	2	$\varepsilon[B s_1 IBI]\varepsilon$	2	$\vec{B} s_1 \vec{IB}$	2
3	2	$s_1 I s_1 I R$	$\varepsilon[BI s_1 IBII]B\varepsilon$	3	$\varepsilon[BI s_2 IBI]\varepsilon$	3	$\vec{BI} s_2 \vec{IB}$	3
4	3	$s_1 B s_2 B R$	$\varepsilon[BII s_1 BII]B\varepsilon$	4	$\varepsilon[BI s_2 BI]\varepsilon$	4	$\vec{BI} s_2 \vec{BIB}$	4
5	4	$s_2 I s_2 I R$	$\varepsilon[B IIB s_2 II]B\varepsilon$	5	$\varepsilon[BIB s_2 II]\varepsilon$	5	$\vec{BIB} s_2 \vec{IB}$	5
6	4	$s_2 I s_2 I R$	$\varepsilon[B IIBI s_2 I]B\varepsilon$	6	$\varepsilon[BIBI s_3 I]\varepsilon$	6	$\vec{BI} s_3 \vec{IB}$	6
7	5	$s_2 B s_3 B L$	$\varepsilon[B IIBII s_2 B]\varepsilon$	7	$\varepsilon[BIBI s_3 B]\varepsilon$	7	$\vec{BI} s_3 \vec{B}$	7
8	6	$s_3 I s_4 B L$	$\varepsilon[B IIBI s_3 IB]\varepsilon$	8	$\varepsilon[BI s_3 IB]\varepsilon$	8	$\vec{BI} s_3 \vec{IB}$	8
9	8	$s_4 I s_5 I L$	$\varepsilon[B IIB s_4 IBB]\varepsilon$	9	$\varepsilon[BIB s_4 IB]\varepsilon$	9	$\vec{BIB} s_4 \vec{IB}$	9
10	10	$s_5 B s_6 B L$	$\varepsilon[BII s_5 BIBB]\varepsilon$	10	$\varepsilon[BI s_5 BIB]\varepsilon$	10	$\vec{BI} s_5 \vec{BIB}$	10
11	11	$s_6 I s_6 I L$	$\varepsilon[BI s_6 IBIBB]\varepsilon$	11	$\varepsilon[BI s_6 IBIB]\varepsilon$	11	$\vec{BI} s_6 \vec{IB}$	11
12	11	$s_6 I s_6 I L$	$\varepsilon[B s_6 IIBIBB]\varepsilon$	12	$\varepsilon[B s_6 BBIB]\varepsilon$	12	$\vec{B} s_6 \vec{BIB}$	12
13	12	$s_6 B s_0 B R$	$\varepsilon[B s_6 B IIBIBB]\varepsilon$	13	$\varepsilon[B s_6 BIBIB]\varepsilon$	13	$\vec{B} s_6 \vec{BIB}$	12
14	1	$s_0 I s_1 B R$	$\varepsilon[BB s_0 IIBIB]\varepsilon$	1'	$\varepsilon[B s_0 IBIB]\varepsilon$	1	$\vec{B} s_0 \vec{IB}$	1
15	2	$s_1 I s_1 I R$	$\varepsilon[BBB s_1 IBIBB]\varepsilon$	2'	$\varepsilon[B s_1 IBIB]\varepsilon$	2'	$\vec{B} s_1 \vec{IB}$	2
16	3	$s_1 B s_2 B R$	$\varepsilon[BBBI s_1 BIBB]\varepsilon$	4'	$\varepsilon[BI s_1 BIB]\varepsilon$	14	$\vec{BI} s_1 \vec{BIB}$	3
17	4	$s_2 I s_2 I R$	$\varepsilon[BBBIB s_2 IBB]\varepsilon$	5'	$\varepsilon[BIB s_2 IB]\varepsilon$	5'	$\vec{BIB} s_2 \vec{IB}$	5
18	5	$s_2 B s_3 B L$	$\varepsilon[BBBIBI s_2 BB]\varepsilon$	7'	$\varepsilon[BIBI s_2 B]\varepsilon$	15	$\vec{BI} s_2 \vec{B}$	13
19	6	$s_3 I s_4 B L$	$\varepsilon[BBBIB s_3 IBB]\varepsilon$	8'	$\varepsilon[BIB s_3 IB]\varepsilon$	16	$\vec{BIB} s_3 \vec{IB}$	14
20	7	$s_4 B s_H B L$	$\varepsilon[BBBI s_4 BBBB]\varepsilon$	14	$\varepsilon[BI s_4 B]\varepsilon$	17	$\vec{BI} s_4 \vec{B}$	15
21	x	$(s_{HALT}, I) \notin D_{\vec{\eta}}$	$\varepsilon[BBB s_{HALT} IB BBB]\varepsilon$	15	$\varepsilon[B s_{HALT} IB]\varepsilon$	18	$\vec{B} s_{HALT} \vec{IB}$	16
xx	xx	$xxxxx$	$xxxxx$	xx	$xxxxx$	xx	$xxxxx$	xx

Table 4. Tracing and staging for Example IV

7. Conclusion

The *unsolvable decision problems* are of two types. The first type of the problem is solvable but not with the objects and decision-counting methods we have at our disposal. The example is the unsolvability of the binomical equations in the real axis. But with the Complex Numbers Theory they are solvable describing the physical reality. The help is that the *imaginary axis* (the new dimension) has been introduced. The another example is the Great Fermat Theorem and its solution (Andrew Wiles 1993 and Ann. Math. 1995).

The second type of the unsolvable or *undecidable* problems are those which are given mistakenly by *having an Auto-Reference embedded*. They are the *paradoxes*, which invokes the infinite cycles when they are 'solved' and just for this they are reducible to the Halting Problem; their

solution without a certain 'step-aside' requires the Perpetuum Mobile functionality. This doesn't mean that a counting under their description is not performed physically but is not resultative in a finite time.

(Nevertheless we can want to have the *infinite cycle for technology purposes*, e.g. the *push-pull circuit*. Here the infinite cycle's functionality is *created intentionally* and, as such, the push-pull circuit is the example of the *recurring but not recursive counting*. Here the Auto-Reference is introduced intentionally, as a successive figuring method, creating an infinite sequence of wanted values.)

Generally, a sequence of states or figuring steps of solving a problem could be *divergent* or *constant* or *convergent*. The divergent and constant cases are felt as the *real example of the infinite cycle in the very sense of this term*. We can say, rather jokingly, that the convergent counting halts, even if it was in the infinity (e.g. the Newton method) and that the divergent counting doesn't halt even in the infinity, including now the constant sequence too - the model is the information transfer in an interrupted information transfer channel. When in the recurring counting the number of figuring steps is not given explicitly, then, the results from the successive steps must be compared. When it is set badly or is not set at all, the infinite cycle occurs and, by the algorithm's definition requiring resultativeness, such a counting is erroneous. The flagrant example of the badly set task is the way in which the Gibbs Paradox arises - here it is the Auto-Reference embedded by not respecting the difference between what is measured (observed), and what is measuring (observing). We used it in the extreme case (with the complete mixing these two levels) as the physical model of all noetic paradoxes.

The aim of this paper was to detect the infinite cycle from its own characteristics. Our vision is that the *counting itself is of the physical character* and, as such, is *subjected to physical laws*, especially, to the II Principle of Thermodynamics. The infinite cycle is viewed as a certain type of an equilibrium state.¹⁷ To await the finite-time end of such states is a paradoxical and, as such, unachievable wish. Nevertheless, all these cycles are representable by the Carnot Cycle (it is the infinite cycle conceptually) used as the thermodynamic model of a cyclic information transfer [6, 8, 10]. We see here the *growth* of thermodynamic entropy within the *whole isolated system* in which the cycle, or the information transfer, is running and we see the *constant* or *decreasing* thermodynamic entropy within its *working medium*, or within the transfer channel in the information-thermodynamic representation. From this point of view the aim to recognize any infinite cycle, to decide the Halting Problem, is solvable. The information-thermodynamic considerations were expressed in terms of the Automaton Theory. The general configuration types of the observed Turing Machine were generated and the Pumping Lemma was used. The author believes that he has shown that problems given paradoxically, erroneously as for being resultant, have the Auto-Reference embedded both *in* the sense of the *objective* of the problem or *in* the sense of *the solving* the problem - the *Auto-Reference can be in the solving method while the very objective of the problem can be solvable*. The author's wish is that

¹⁷ The interesting is that the stability of an equilibrium state and of an atomic structure are similar. Without the natural radioactivity the end of atoms seems to 'be in the infinity', too.

the following claim could be considered as the theorem for recognizing, deciding of any infinite cycle:

Due to the fact that any infinite cycle starts at a finite time and, for the Control Unit of any Turing Machine is an finite-state automaton, and due to the fact that the Pumping Lemma it is valid for the regular infinite and thus periodical language of the general configuration types of the observed Turing Machine, the Halting Problem is decidable.

8. Appendix

We consider the basic types of chains of the terminal symbols on the input-output TM 's tape,

$$\mathbf{I} \triangleq \vec{I}I = I, \vec{I} = II, \vec{I} = III \dots$$

$$\mathbf{B} \triangleq \vec{B}B = B, \vec{B} = BB, \vec{B} = BBB, \underline{\mathbf{B}} \triangleq \varepsilon$$

Further types are

$$\mathbf{IB} \triangleq \vec{I}\vec{B}; \vec{I}\vec{B} = IB, \vec{I}\vec{B} = IIB, \vec{I}\vec{B} = III \dots B$$

$$\vec{I}\vec{B} = IBB, \vec{I}\vec{B} = IBBB \dots$$

$$\vec{I}\vec{B} = III \dots BBB \dots$$

$$\mathbf{BI} \triangleq \vec{B}\vec{I}; \vec{B}\vec{I} = BI, \vec{B}\vec{I} = BBI, \vec{B}\vec{I} = BBB \dots I$$

$$\vec{B}\vec{I} = BII, \vec{B}\vec{I} = BIII \dots$$

$$\vec{B}\vec{I} = BBB \dots III \dots$$

$$\mathbf{IBI} \triangleq \vec{I}\vec{B}\vec{I} = III \dots BBB \dots III \dots I$$

$$\mathbf{BIB} \triangleq \vec{B}\vec{I}\vec{B} = BBB \dots III \dots BBB \dots B$$

$$\mathbf{IBB} \triangleq \mathbf{IB}$$

$$\mathbf{IBB} \triangleq \mathbf{IB},$$

$$\mathbf{IIB} \triangleq \mathbf{IB}$$

$$\mathbf{IIB} \triangleq \mathbf{IB}$$

$$\vec{\mathbf{IB}} \triangleq \vec{\mathbf{IB}}, \vec{\mathbf{IB}} \triangleq \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBIB}}$$

$$\vec{\mathbf{IBB}} \triangleq \vec{\mathbf{IBB}}, \vec{\mathbf{IBB}} \triangleq \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBB}} = \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBIB}} = \vec{\mathbf{IB}}$$

The following equivalences of chains of the terminal symbols are considerable,

$$\vec{\mathbf{IBB}} \triangleq \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBIBB}} = \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBIB}} = \vec{\mathbf{IB}}$$

$$\vec{\mathbf{IBI}} \triangleq \vec{\mathbf{IBIIBI}} \dots \vec{\mathbf{IBIIBI}} = \vec{\mathbf{IBIB}} \dots \vec{\mathbf{IBI}}$$

$$\vec{\mathbf{IBII}} \triangleq \vec{\mathbf{IBI}}$$

$$\vec{\mathbf{IBIB}} \triangleq \vec{\mathbf{IB}}$$

$$\vec{IBIB} \triangleq \vec{IBIB}, \vec{IBIB} \triangleq \vec{IBIBI} \dots \vec{BIBIBIBIBIBIBIBI} \dots \vec{IBIBIB} = \vec{IB}$$

$$\vec{BBI} \triangleq \vec{BI}$$

$$\vec{BI} \triangleq \vec{BI}$$

$$\vec{BI} \triangleq \vec{BIBI} \dots \vec{BIBI}$$

$$\vec{BII} \triangleq \vec{BII}, \vec{BII} \triangleq \vec{BIBI} \dots \vec{BII} = \vec{BIBI} \dots \vec{BIBI} = \vec{BI}$$

$$\vec{BIB} \triangleq \vec{BIBI} \dots \vec{BIBIB} = \vec{BIB}$$

$$\vec{BIB} \triangleq \vec{BIBBIB} \dots \vec{BIBBIB} = \vec{BIBI} \dots \vec{BIB}$$

$$\vec{BIBI} \triangleq \vec{BI}$$

$$\vec{BIBB} \triangleq \vec{BIB}$$

$$\vec{BIBI} \triangleq \vec{BIBIBIBI}, \vec{BIBI} \triangleq \vec{BIBIB} \dots \vec{IBIBIBIBIBIBI} \dots \vec{BIBIBIBI} = \vec{BI}$$

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